

L8

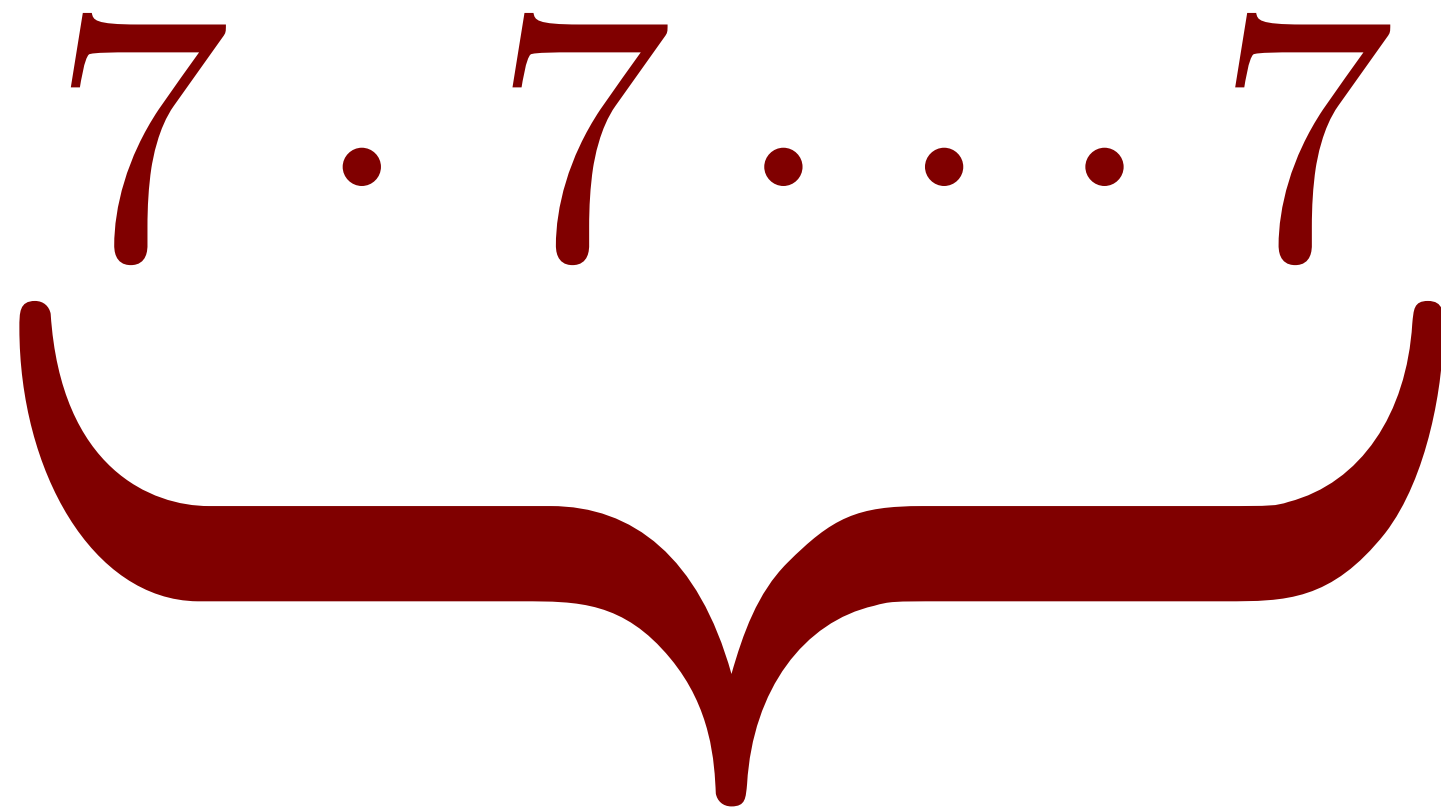
abhi shelat

7

7

*

7



\times times

$$7^x$$

$7^{1000000000000000000}$



[Examples](#) [Random](#)

Input:

$7^{100\,000\,000\,000\,000\,000}$

Power of 10 representation:

$10^{10^{10^{1.228577610529823}}}$

Number length:

84 509 804 001 425 684 decimal digits

$\approx 8.45098 \times 10^{16}$ digits

Last few digits:

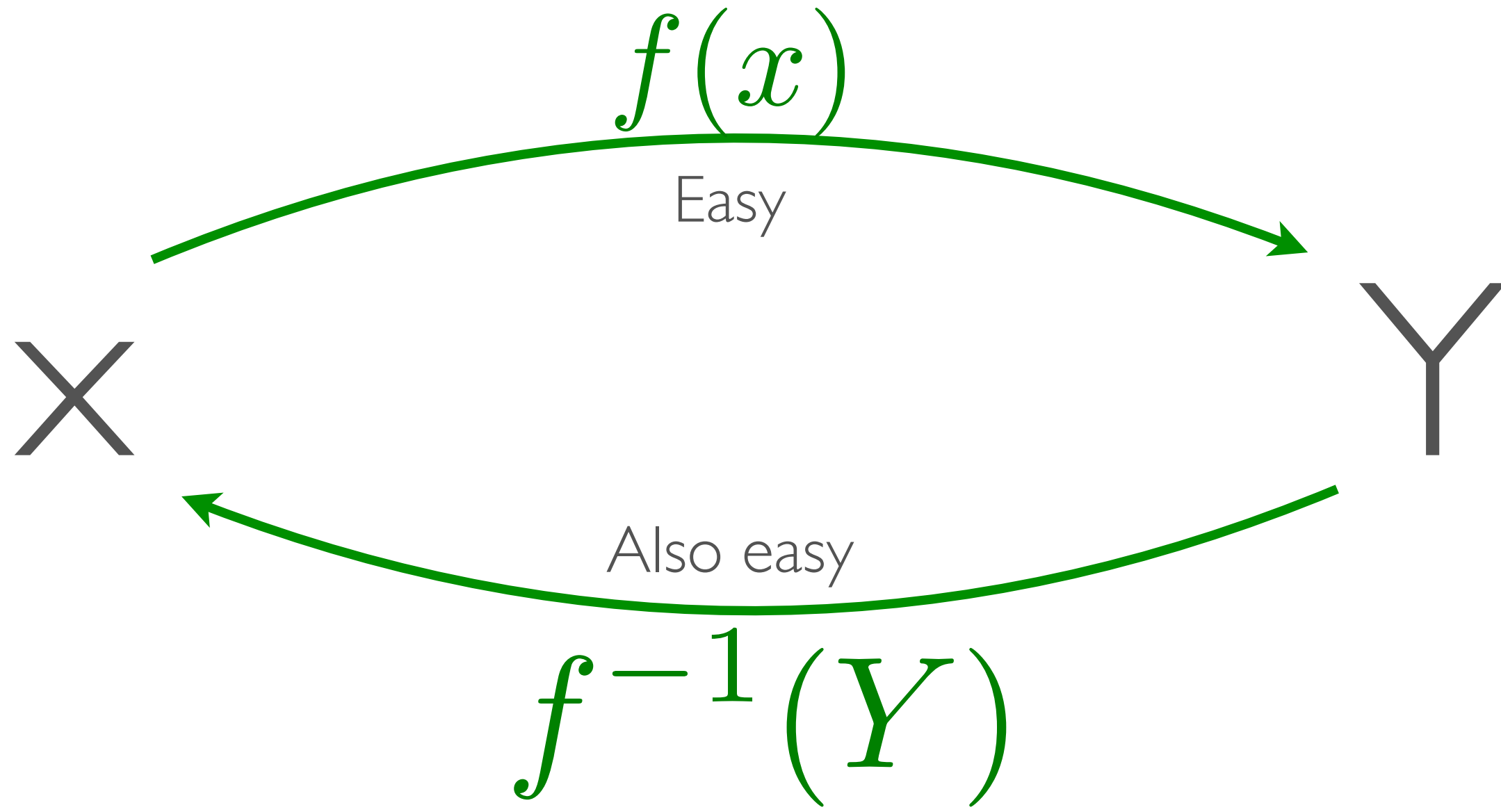
...0000000001

Inverse problem is EASY:

Given: 1231928

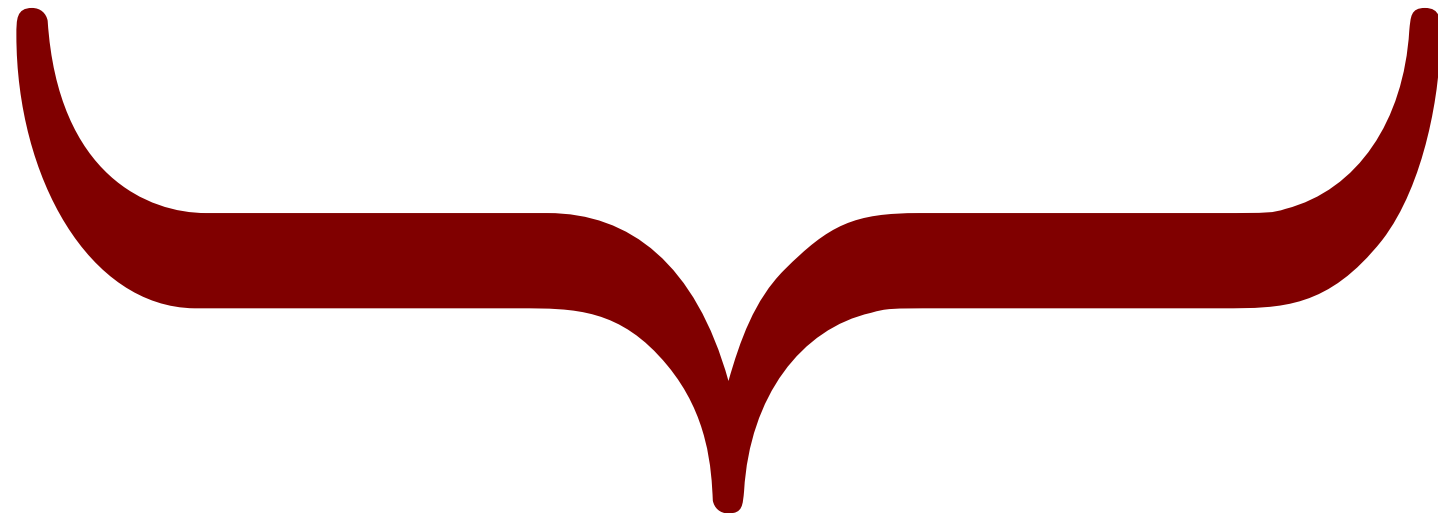
Find X such that $7^X = 1231928$

7.2069571077666163863237014656943296629255654569808791967051



7x

7 · 7 · · · 7



x times

$$7^x \pmod{p}$$

$$7^2 = 7 \cdot 7 \pmod{131}$$

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$$= 49$$

$$7^3 = 49 \cdot 7 \pmod{131}$$

$$343 \pmod{131}$$

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$$\begin{array}{r} 131 \overline{) 343} \\ \underline{262} \\ 81 \end{array} \pmod{131}$$

$$7^3 = 49 \cdot 7 \pmod{131}$$

$$\begin{array}{r} 2 \\ 131 \overline{) 343} \\ \underline{262} \\ 81 \end{array} \pmod{131}$$

$$= 81 \pmod{131}$$

Pick a large prime number, ~200 digits long

$p =$
52061717268811022582985085340962890043057948799610254
54308897584403272194404580619934826278347900750288237
80618996828085538214474401045588718257206979528176365
6923746706192367848073555079554994166440456995548499

$$(7, x, p) \rightarrow 7^x \pmod p$$

Remarkably easy

Assailing U.S. and Kiev, Putin Keeps Open Option of Force, By STEVEN LEE MYERS, ELLEN BARRY and ALAN COWELL 56 minutes ago, President Vladimir V. Putin of Russia on Tuesday described events in Ukraine as an unconstitutional coup, expressed contempt toward the United States and said that any potential use of military force would be a last resort.

sage:

```
power_mod(7,661615980609061103328546834632981100327  
606021844328118170610327602021315328013021032801317  
061210331202327112150001443267213283846986697832766  
969327789698283443269767669783266658282893298110032  
65766578326779876976763253543310061181702153298041  
144328115021606010211163287089801061006143286463281  
181706103312023283181616059733121032851802160098213  
301021600150599020033021902111715330610328608149806  
110132981532981033181100121116170617181706121098083  
300121812443302211315021616020033001211170210131633  
171219981500331705013286110617020032841698170215329  
811003315980600331704981632981121331312170211170598  
083318160133120233100609061698152133031215000133201  
2180900329901329733089816163315021612151646,p)
```

```
497929711688584217527863332300629329098654983269423  
357729595289026203941126793387158661751180509701235  
256633429962635553676002083869590769727522267027725  
835398089223356845936769316025070772103626851899346  
3932429
```

$$Y = 12345$$

$$(7, p, Y) \rightarrow x$$

$$\text{such that } 7^x \bmod p = Y$$

Incredibly hard!

World record in discrete logarithms in $GF(p)$

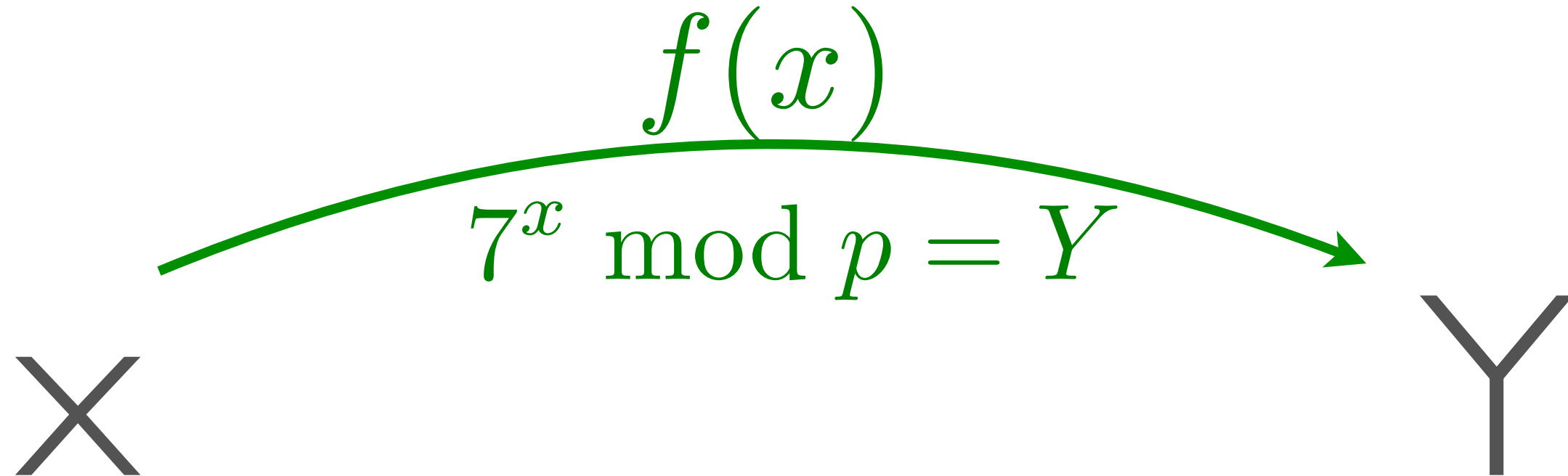
232 digits (768 bits)

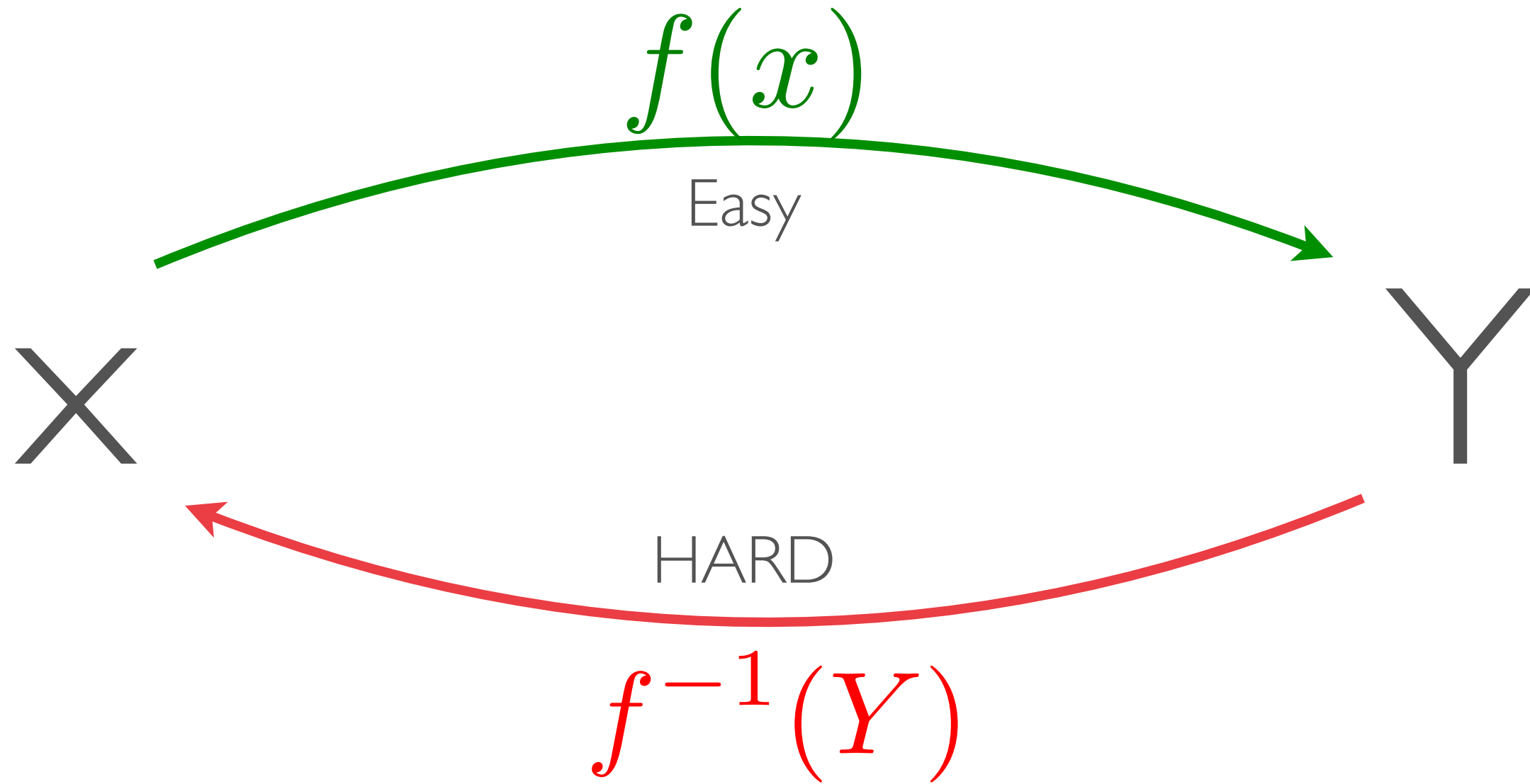
6600yrs of CPU time

Intel Xeon E5-2660 at 2.2 GHz

2016

Thorsten Kleinjung, Claus Diem, [Arjen K. Lenstra](#), Christine Priplata, and Colin Stahlke



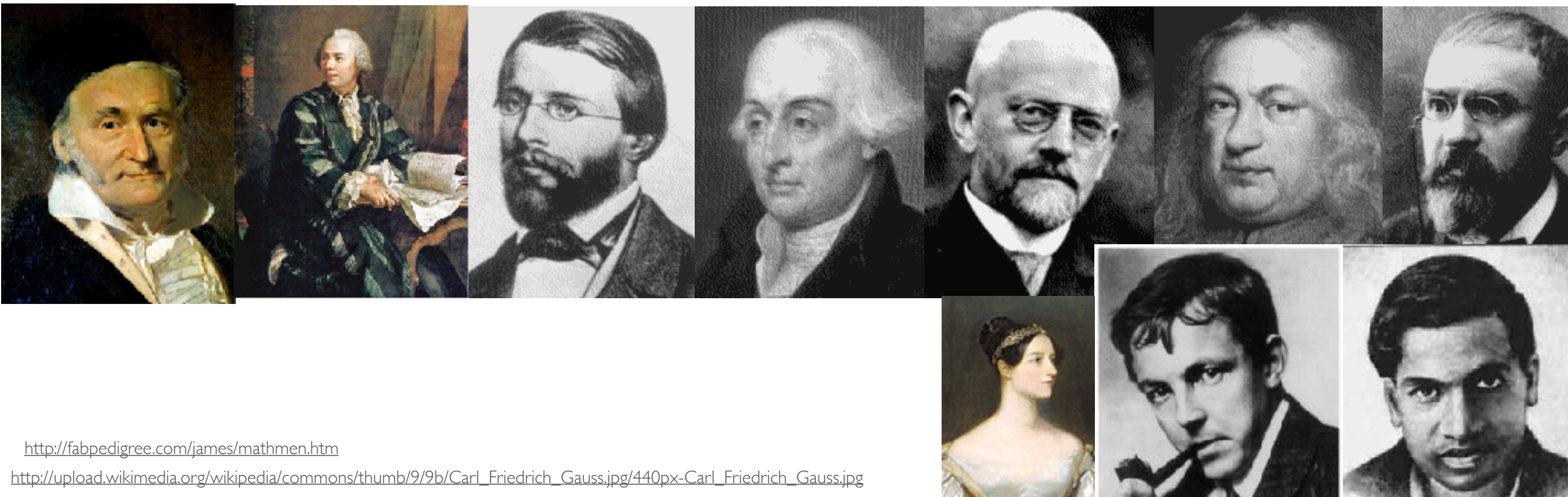
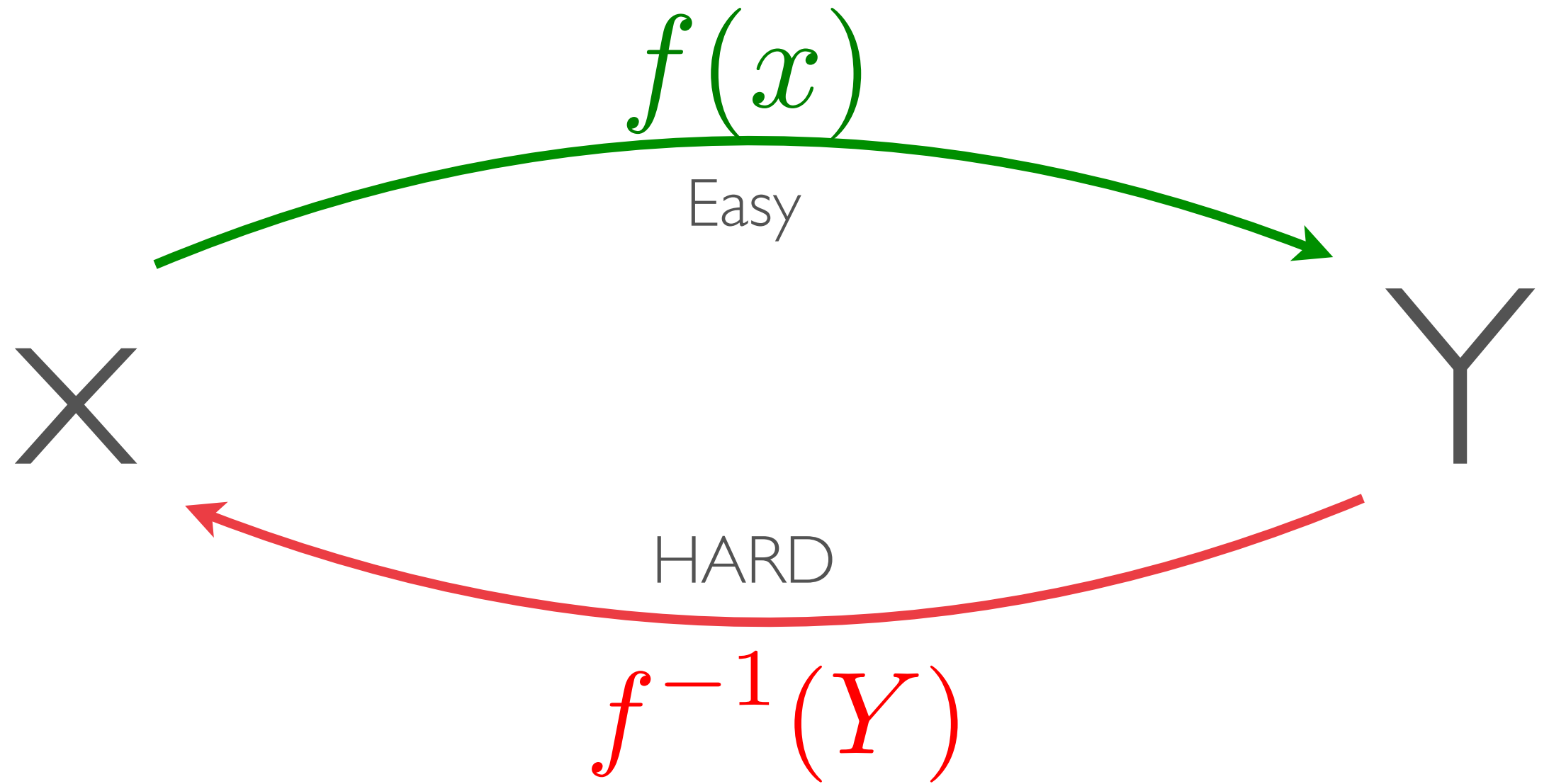


compute an x such that $7^x = Y \pmod p$



<http://www.fitzmuseum.cam.ac.uk/gallery/chinesevases/sortingfragments.html>





<http://fabpedigree.com/james/mathmen.htm>

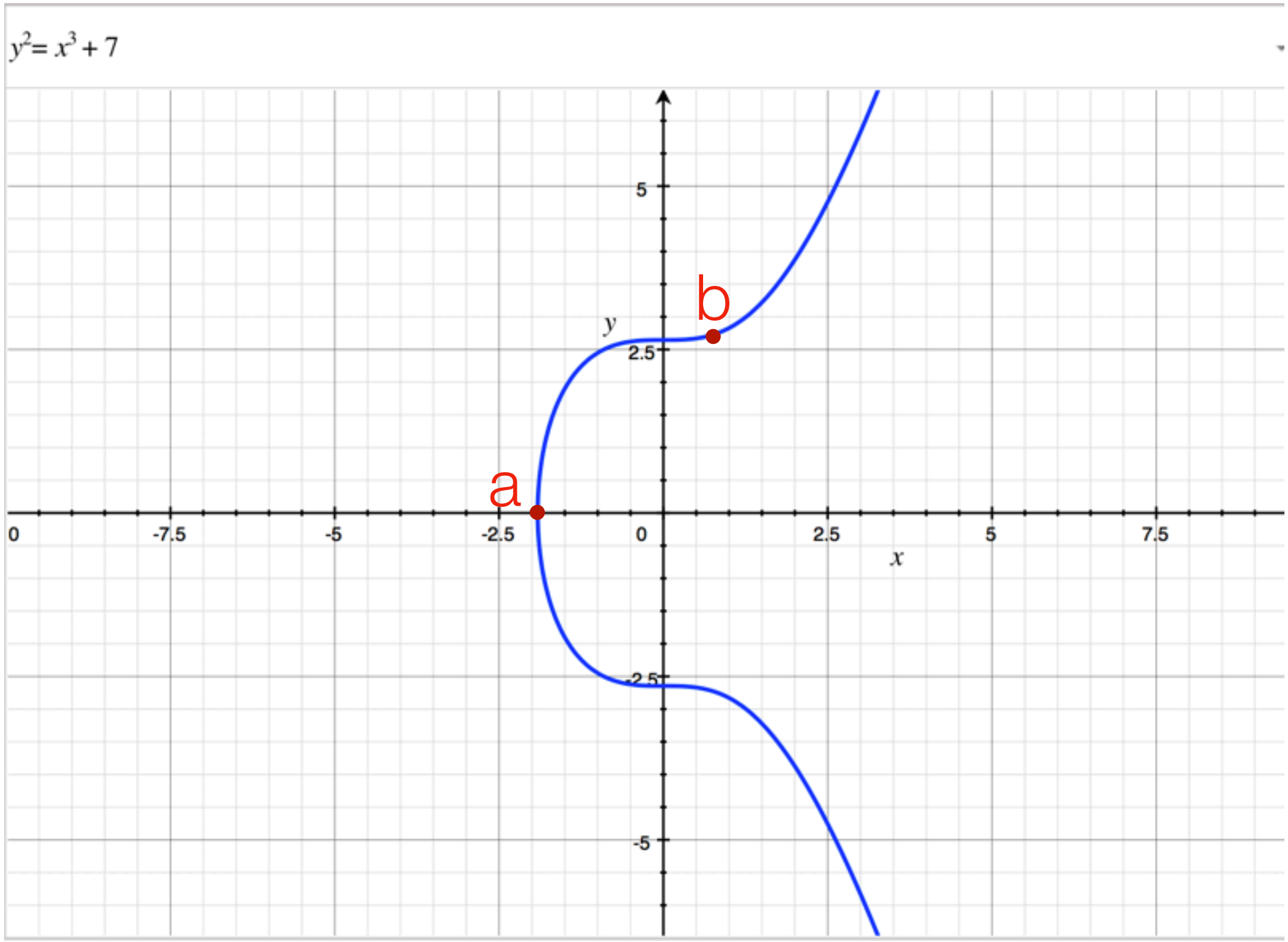
http://upload.wikimedia.org/wikipedia/commons/thumb/9/9b/Carl_Friedrich_Gauss.jpg/440px-Carl_Friedrich_Gauss.jpg

Same problem, harder (seemingly)

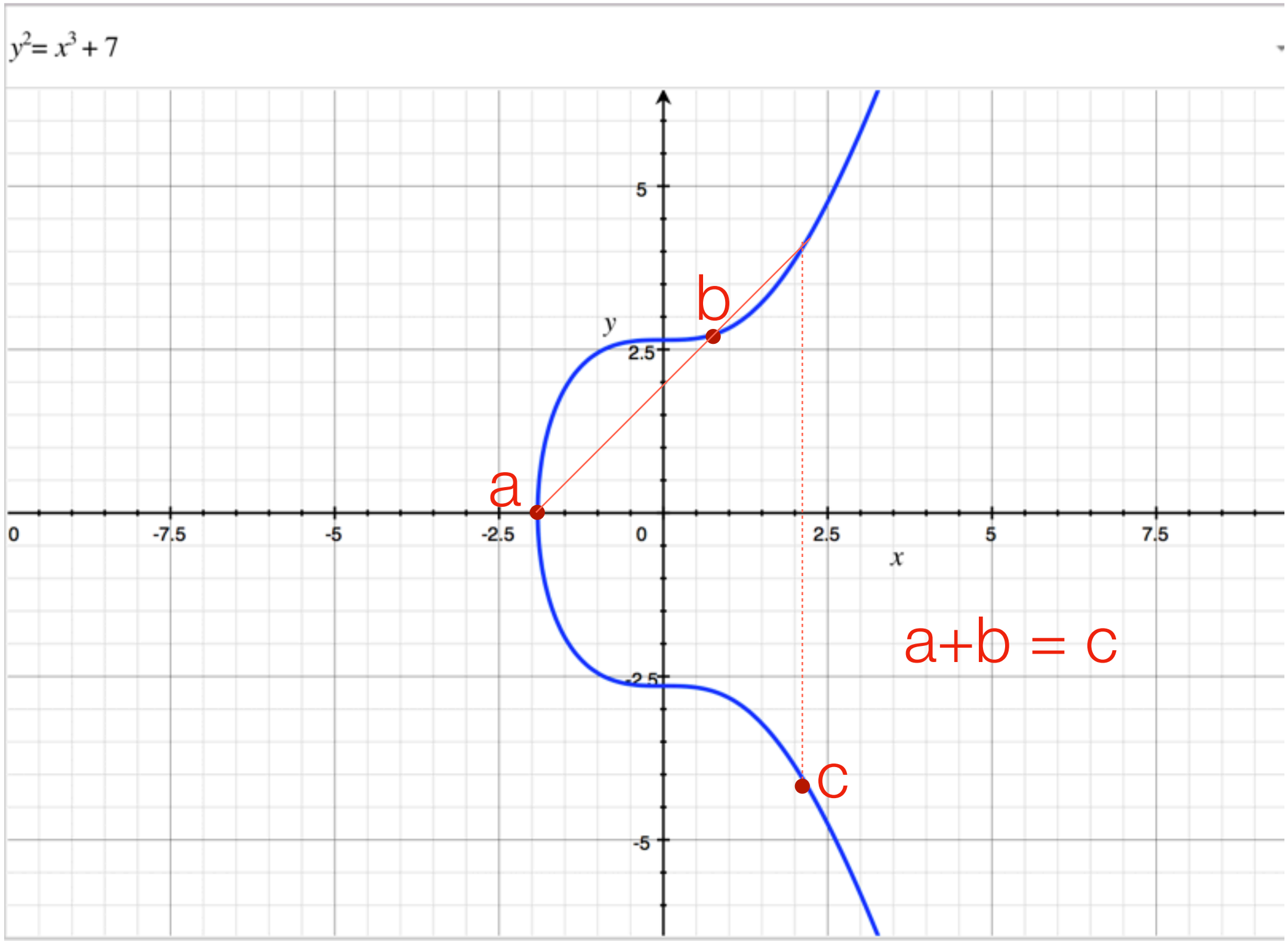
Instead of using 7 (integers mod p), use (x,y) points that lie on an elliptic curve such as

$$y^2 = x^3 + 7$$

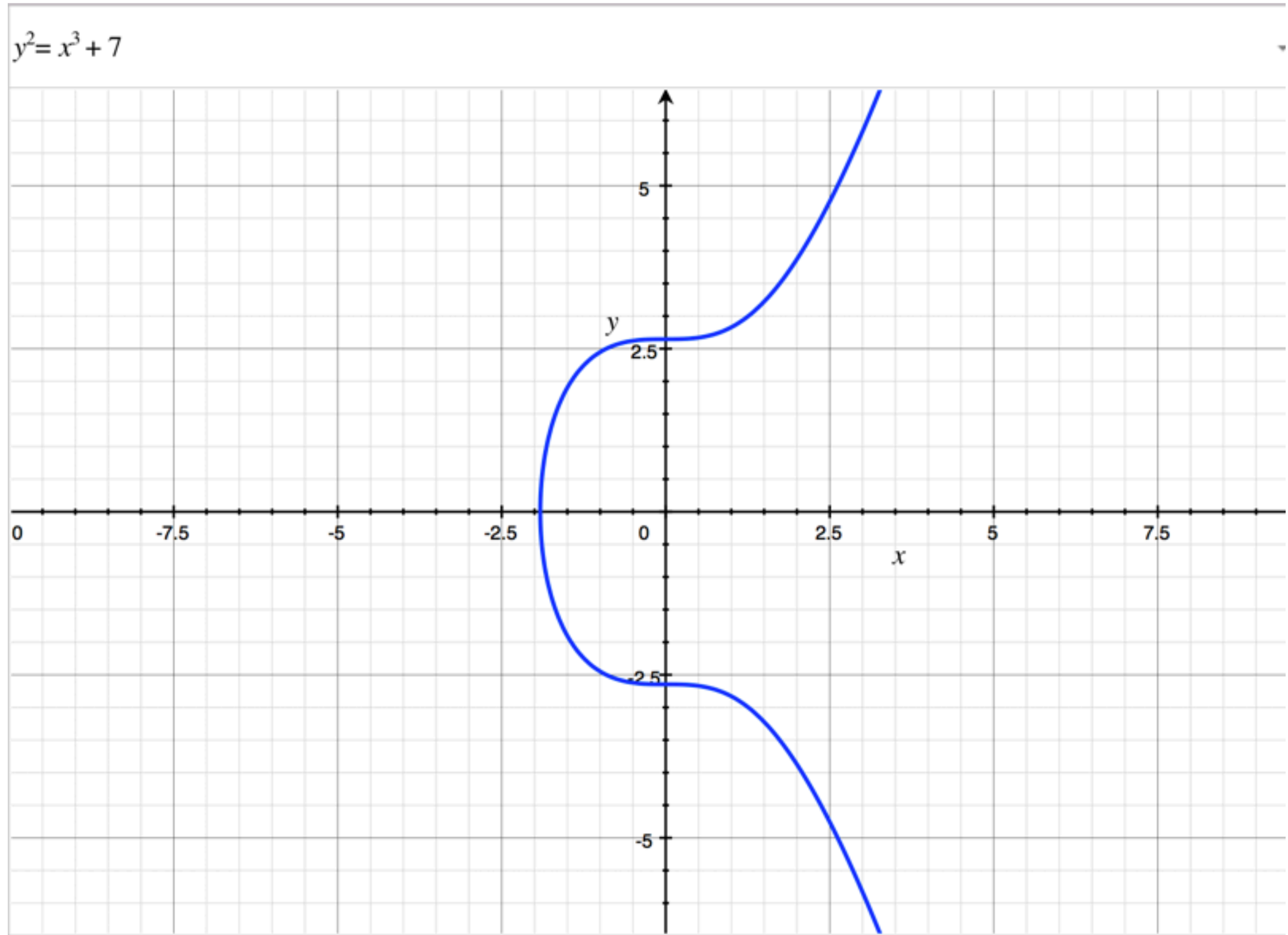
Points on such curves can be added



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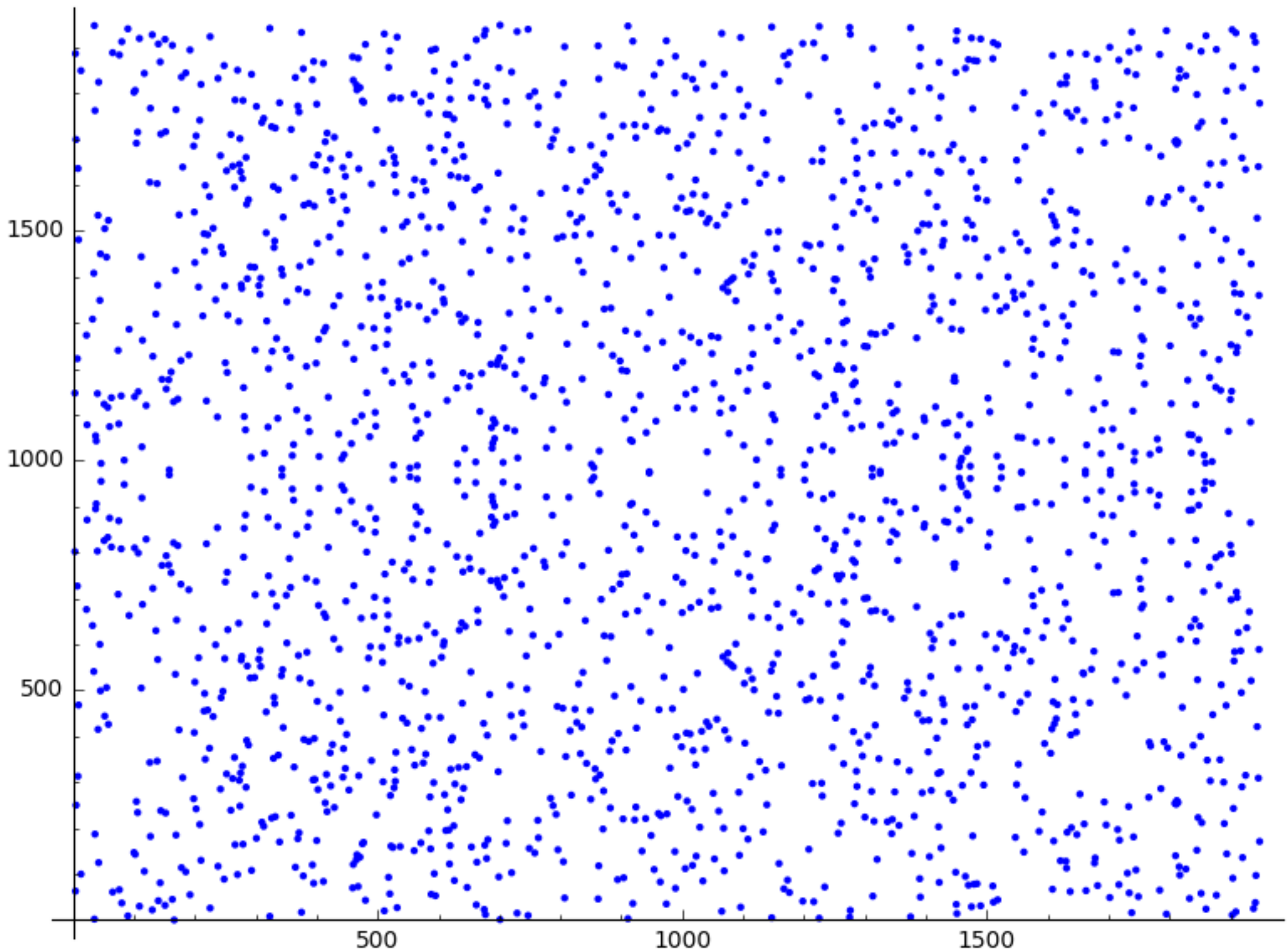


Problem: infinite precision math



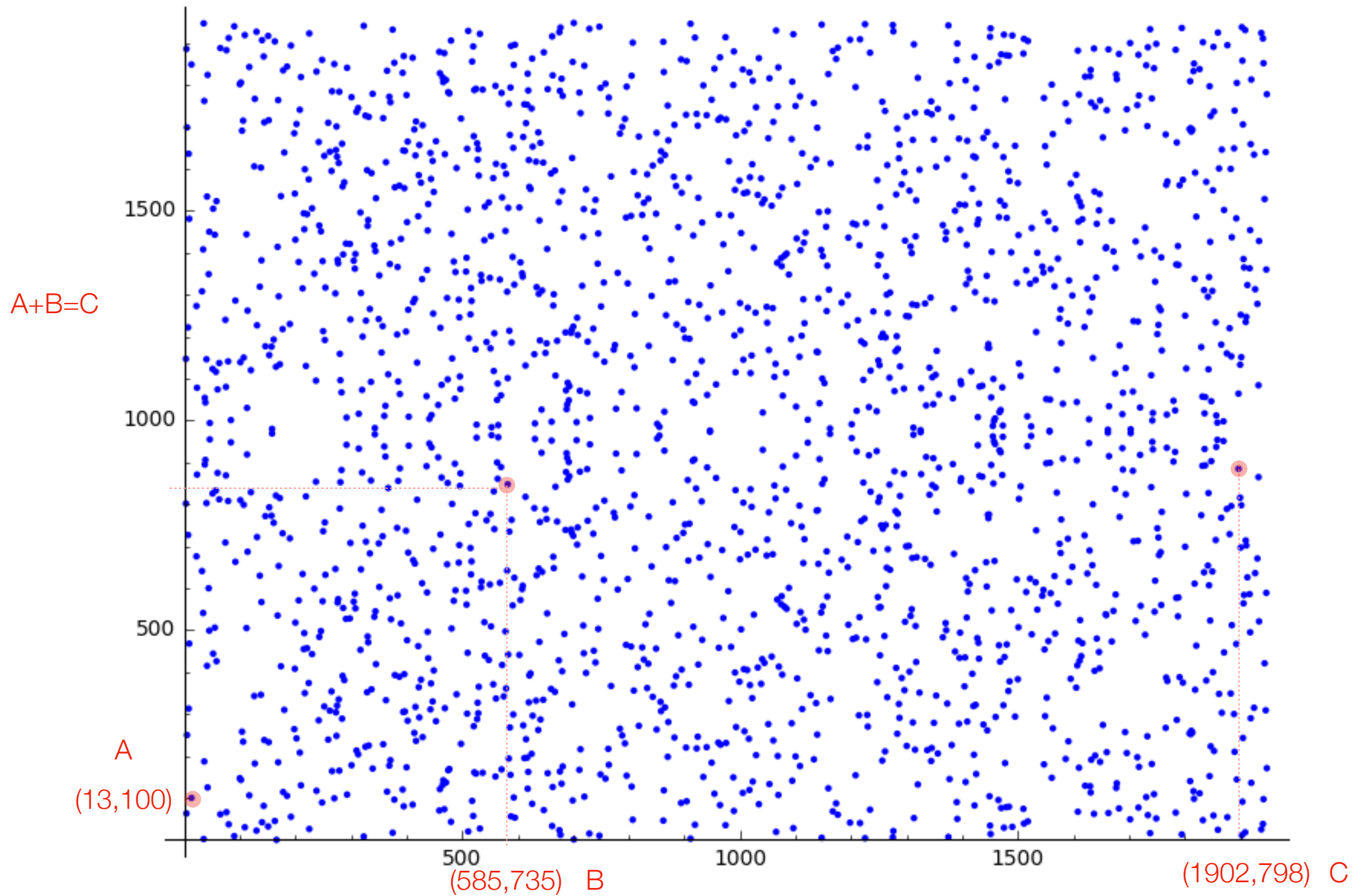
Solution: mod p

$$y^2 = x^3 + 7 \pmod{1949}$$



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Use a big prime

$$y^2 = x^3 + 7 \pmod{115,792,089,237,316,195,423,570,985,008,687,907,853,269,984,665,640,564,039,457,584,007,908,834,671,663}$$

$$2^{256} - 2^{32} - 2^9 - 2^8 - 2^7 - 2^6 - 2^4 - 1$$

The hard problem

Given the points A and $Y = \underbrace{A+A+A+\dots+A}_{x \text{ times}}$

Find x .

This is the **discrete log problem** on elliptic curves

ECDSA signature scheme

We can use this hard problem to make a signature scheme.

One standard way to do this is ECDSA.

This is the sig scheme used in Bitcoin, Ethereum, Litecoin, etc.

ECDSA Gen

Recall $\text{Gen}()$ produces (sk, pk) .

Let G be the elliptic curve and g a **generator** of G .

$\text{Gen}()$: Pick a random x in $[1..q]$

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Output $sk = x$

ECDSA sign

These steps seem magical until you understand how and why they work (in a crypto class).

Sign(sk, msg): $z = \text{hash}(\text{msg})$
 $k = \text{pick random \# in } [1..q-1]$
 $r = \text{“x” coordinate of } k^*g,$
 $s = k^{-1}(z + r^*sk) \text{ mod } q$
Output (r,s) as signature.

If the EC DLOG problem is hard, and hash is secure, this scheme is shown to be secure against forgeries.



abhi — -bash — 125x27

ab2017:abhi abhi\$





abhi — -bash — 125x27

ab2017:abhi abhi\$

