

L9

Hard to agree

abhi shelat

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If party i is not faulty, value $\mathbf{V}_i =$ private value v_i

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4. After receiving (a_1, \dots, a_4) take a majority.

Why does PLS $n=4$ work?

What happens if everyone is honest?

P1 output: (_ , _ , _ , _)

P2 output: (_ , _ , _ , _)

P3 output: (_ , _ , _ , _)

P4 output: (_ , _ , _ , _)

Why does PLS $n=4$ work?

What happens if one party, say P1, is faulty?

P2 output: (__, __, __, __)

P3 output: (__, __, __, __)

P4 output: (__, __, __, __)

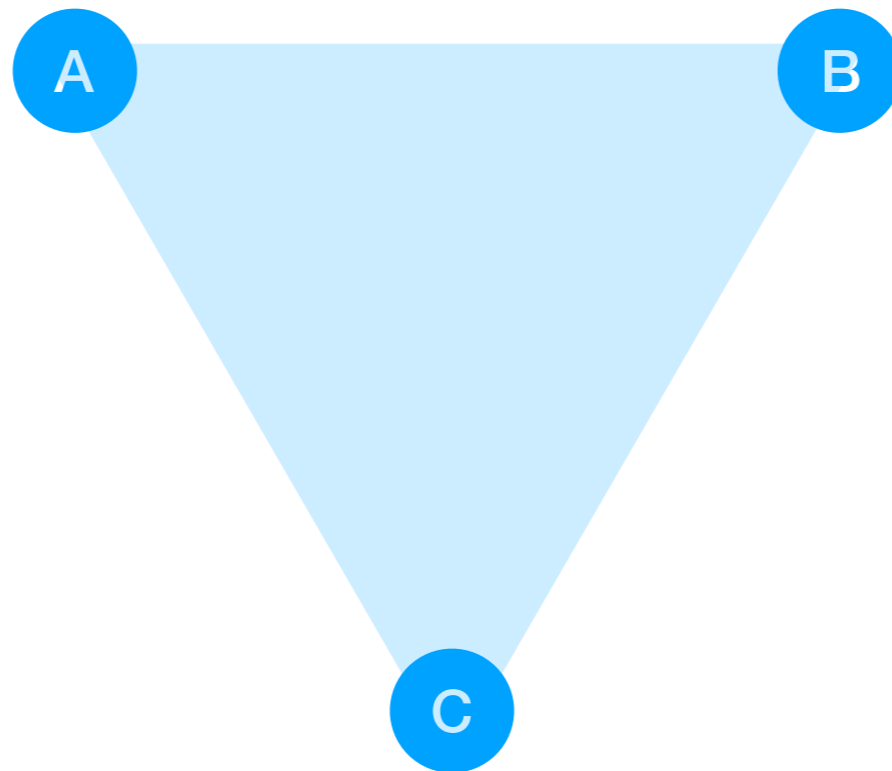
PLS can be generalized

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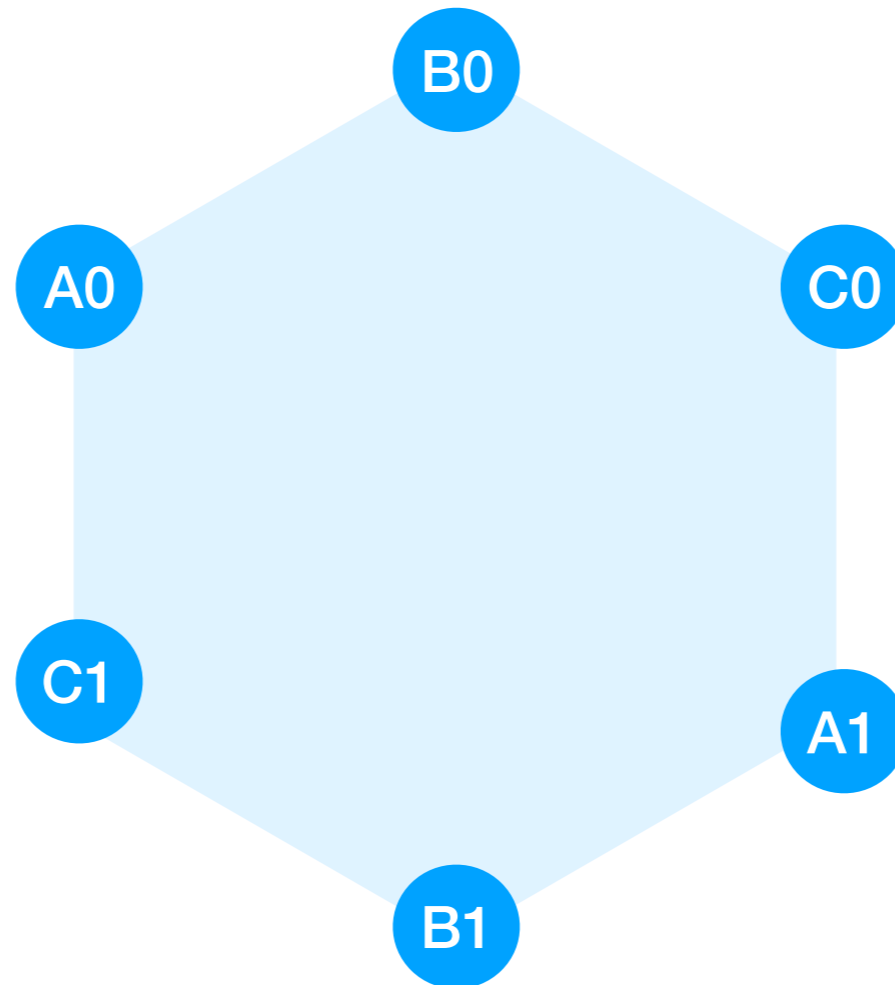
[Fischer-Lynch-Merritt proof]



Consider $n=3$, $m=1$ parties.

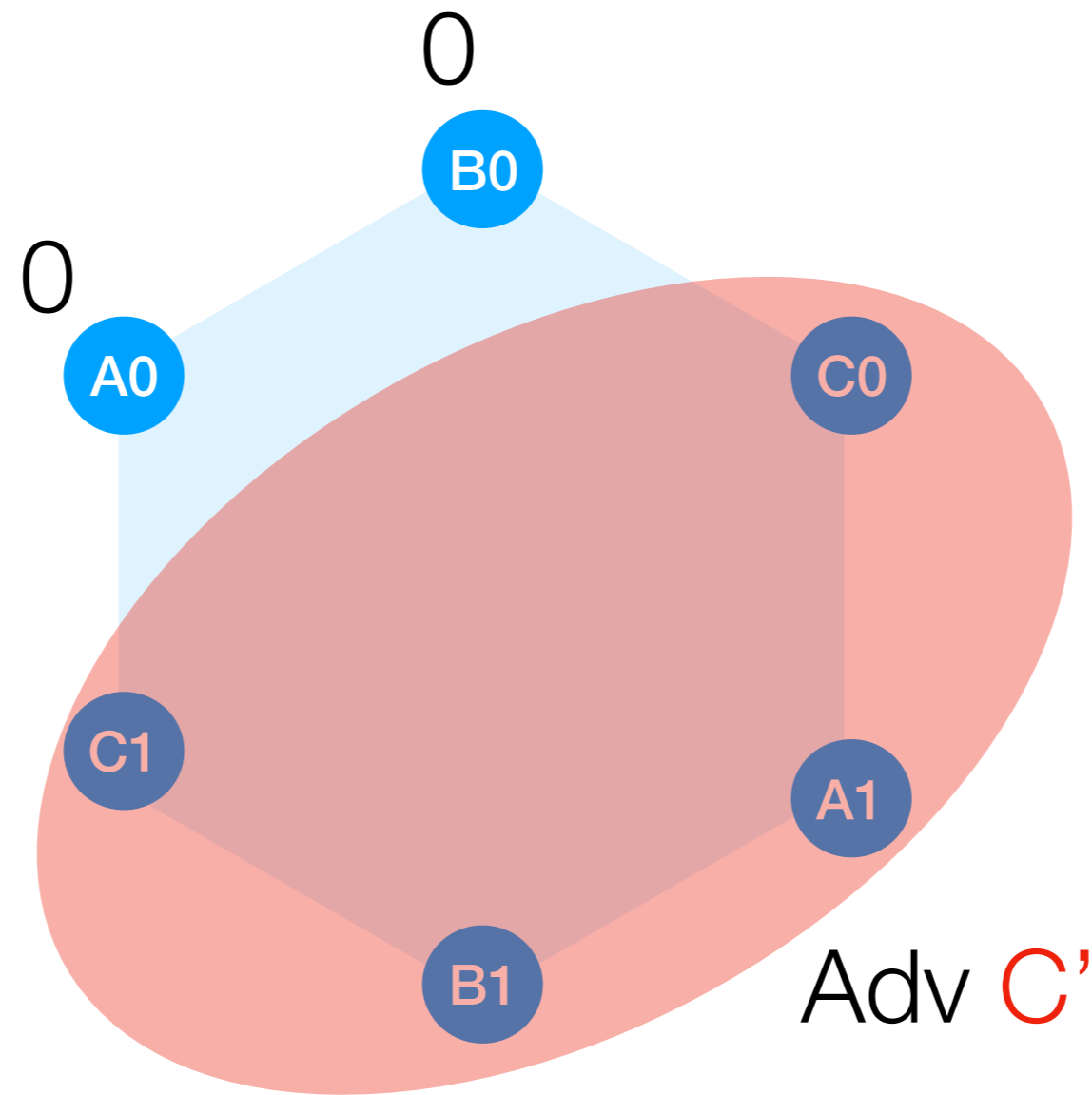
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[Fischer-Lynch-Merritt proof]



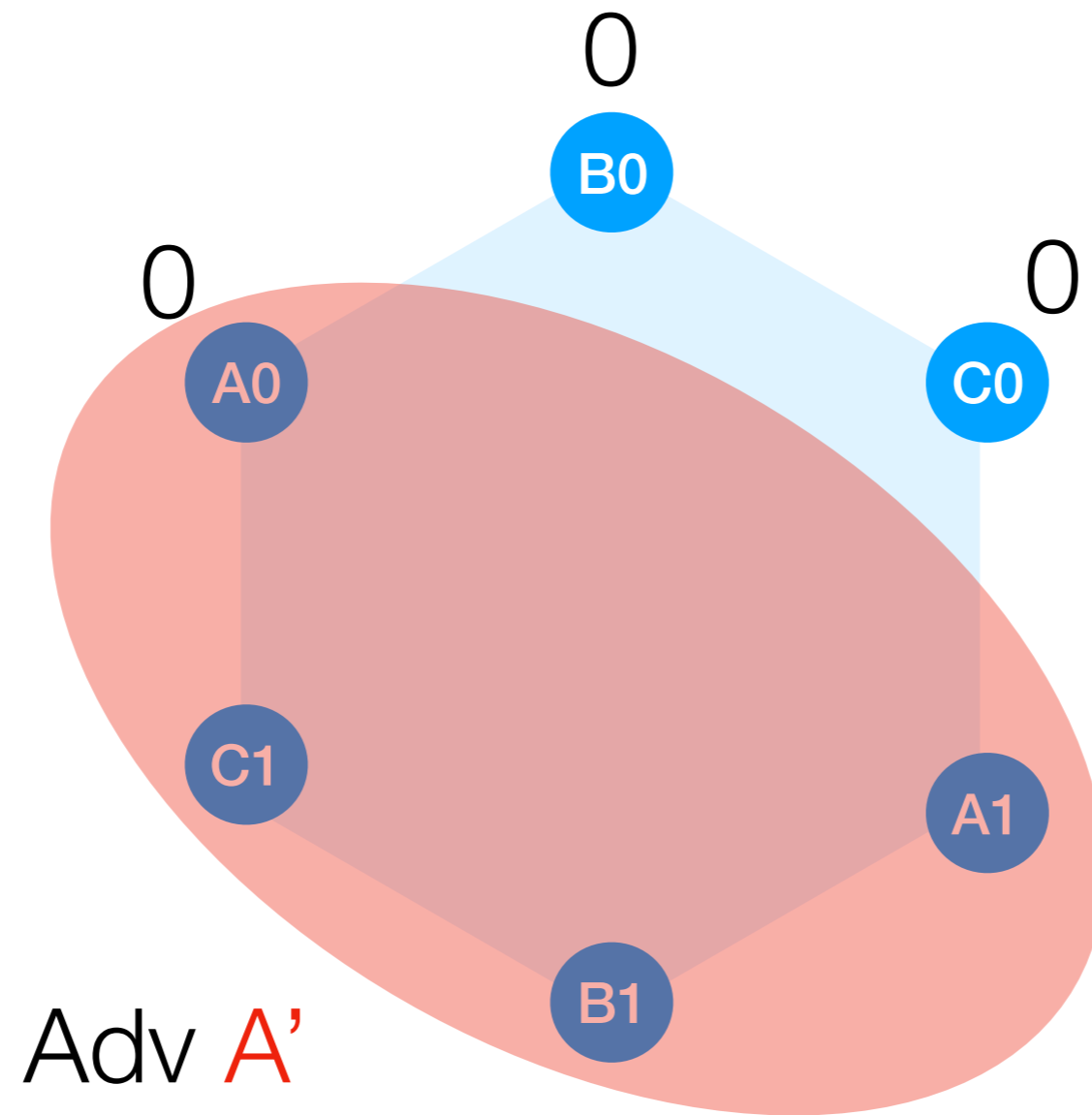
Duplicate parties, P_i running protocol with input i .

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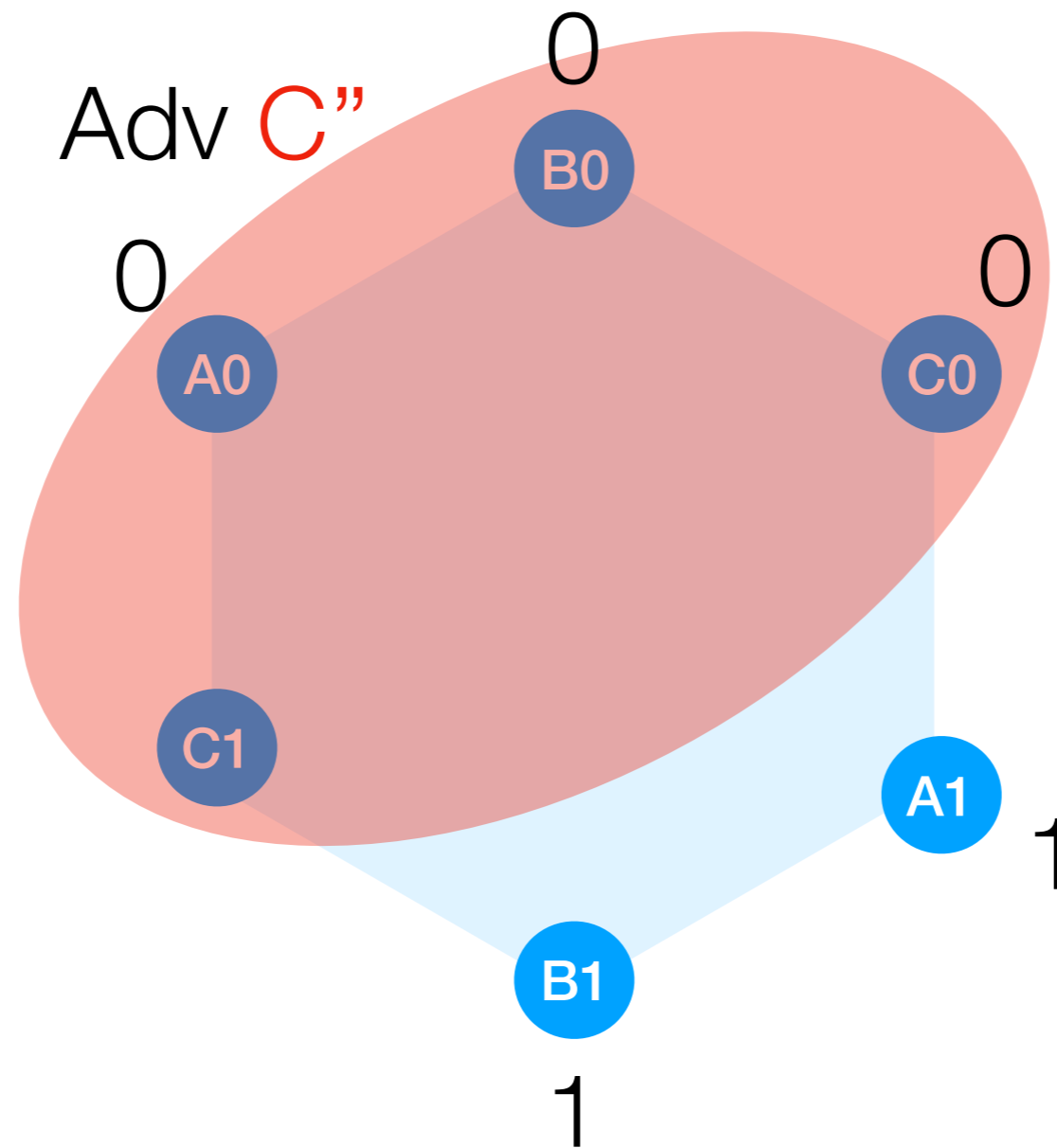
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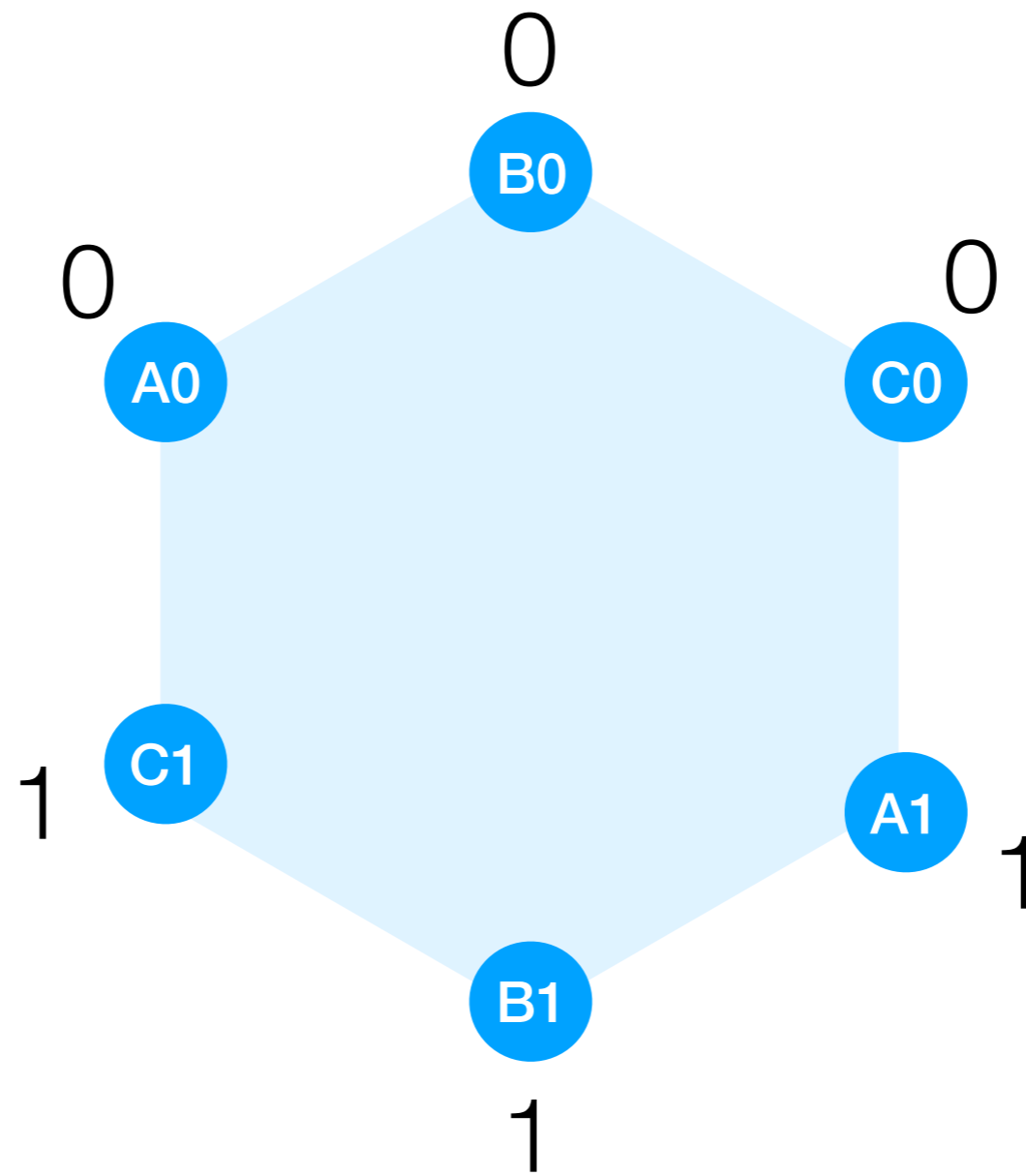
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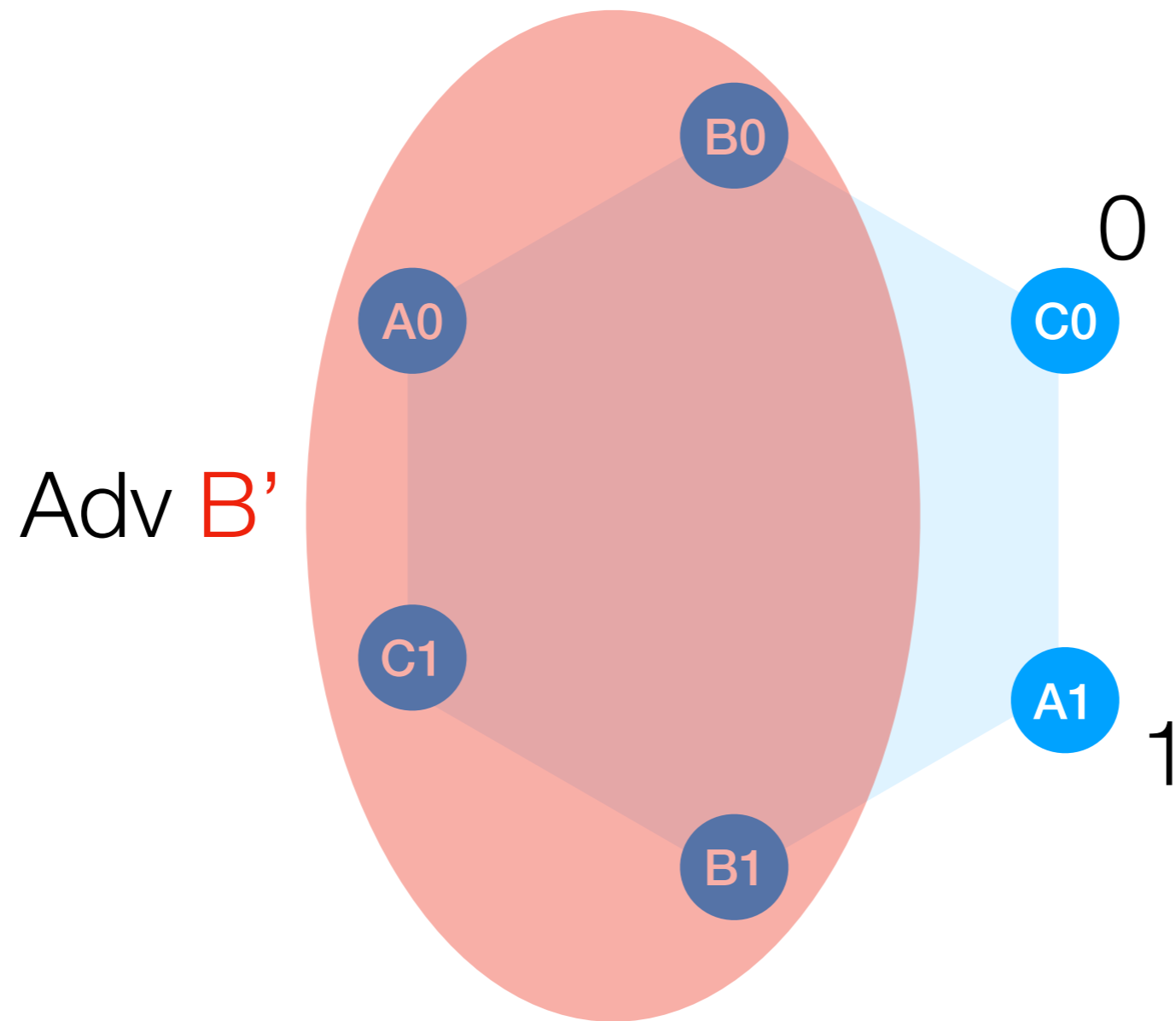
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These must be the outputs.

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Thus, there exists Adv B' which violates consistency.