

L1 5880

Algorithms

abhi shelat

Jan 18 2022

let me intro myself

first goal: create an
amazing learning
experience

second goal: share
basic beautiful ideas
from computer science

third goal: help
prepare you for a job
in CS

what is this
course about?

Theme 1

Small problems are easy to solve

Solve large problems by recasting
them into smaller ones.

Theme 1

Small problems are easy to solve

Solve big problems by making them into smaller ones

Theme 2

Learning how to convince through reason is a great mark of understanding

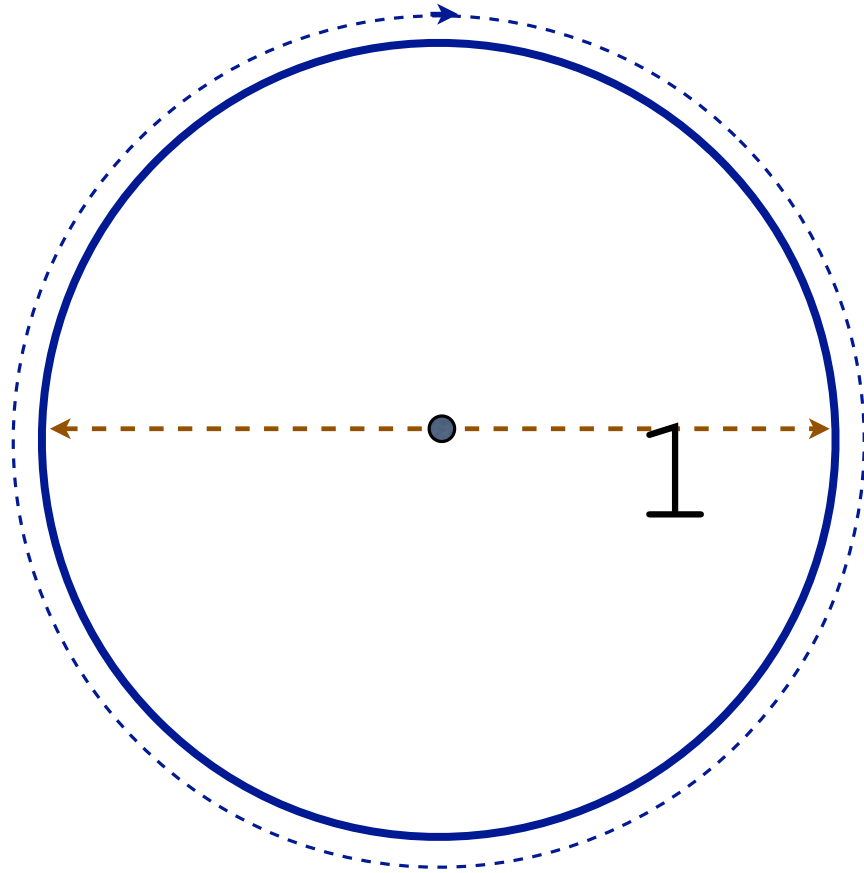


great pyramid at giza 2500bc

image from wikimedia

π





“how much granite/glass do i need?”

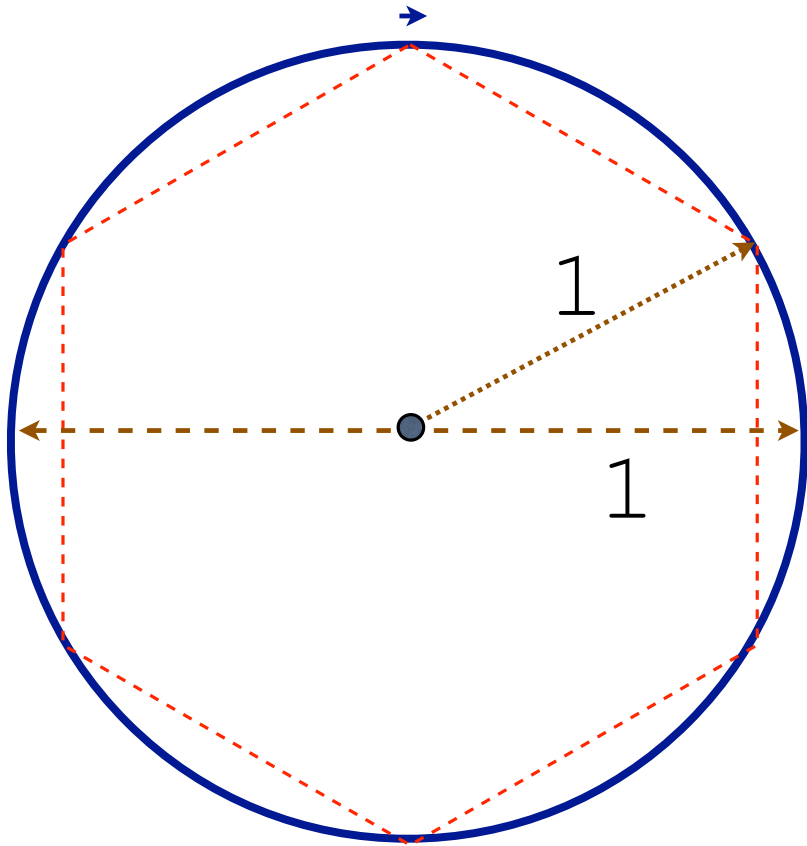
algorithm
to compute

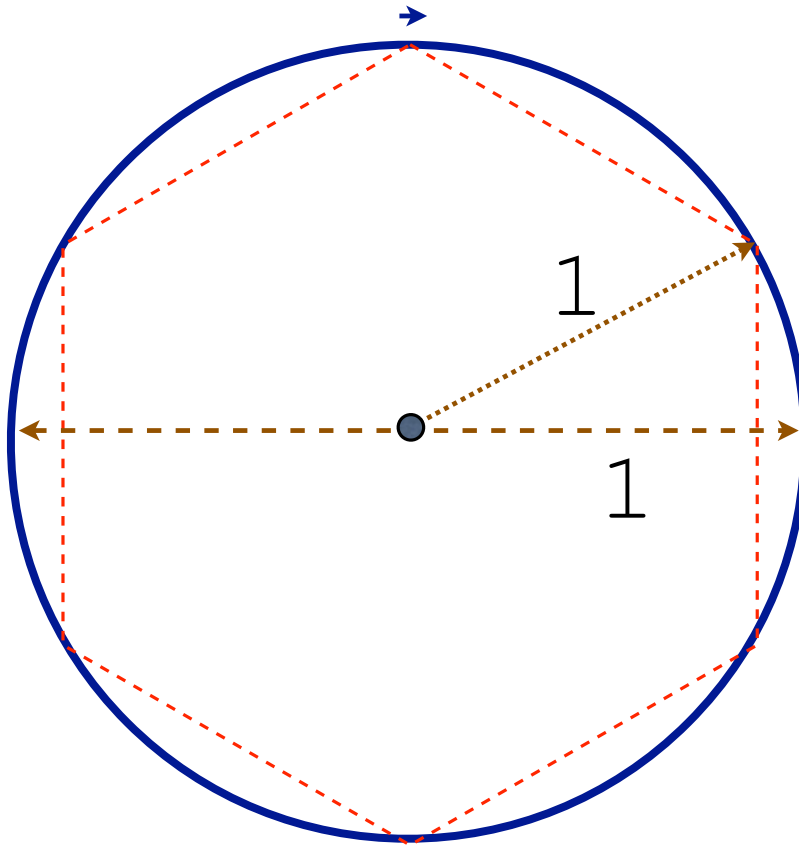
π

250 BCE

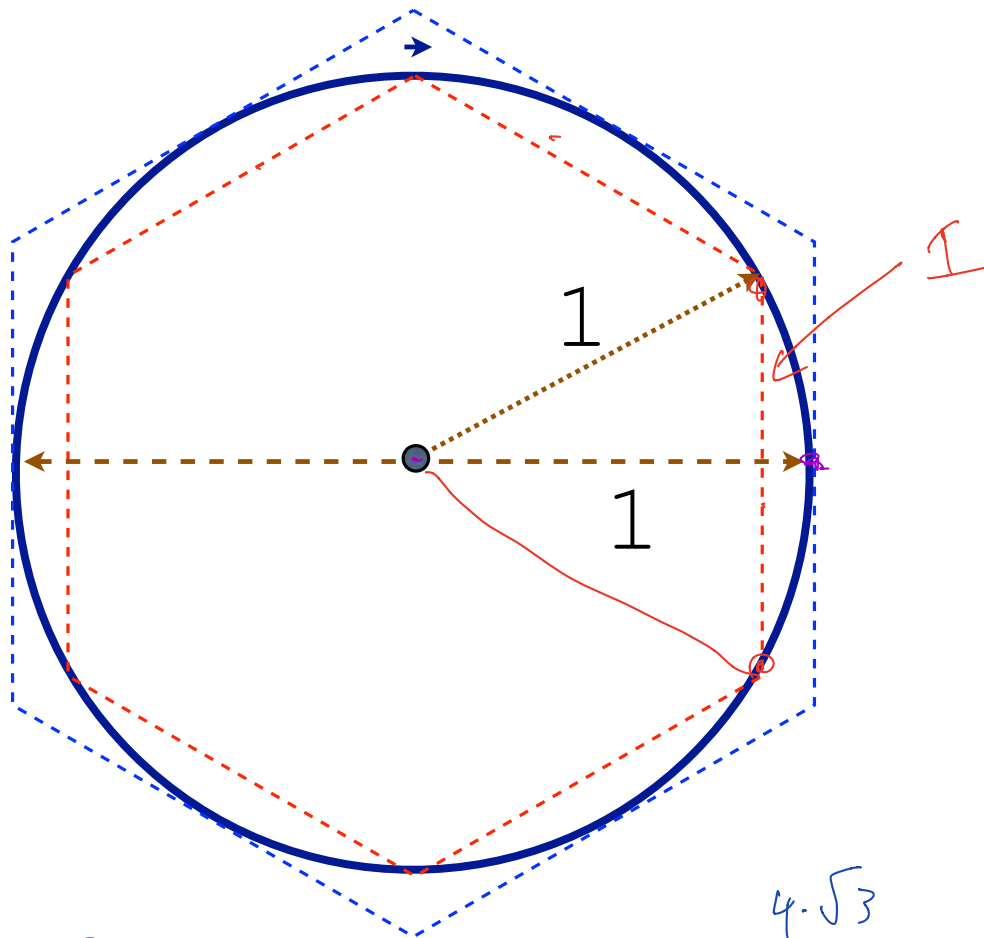
Written down by Archimedes

Idea: It is OK
to approximate

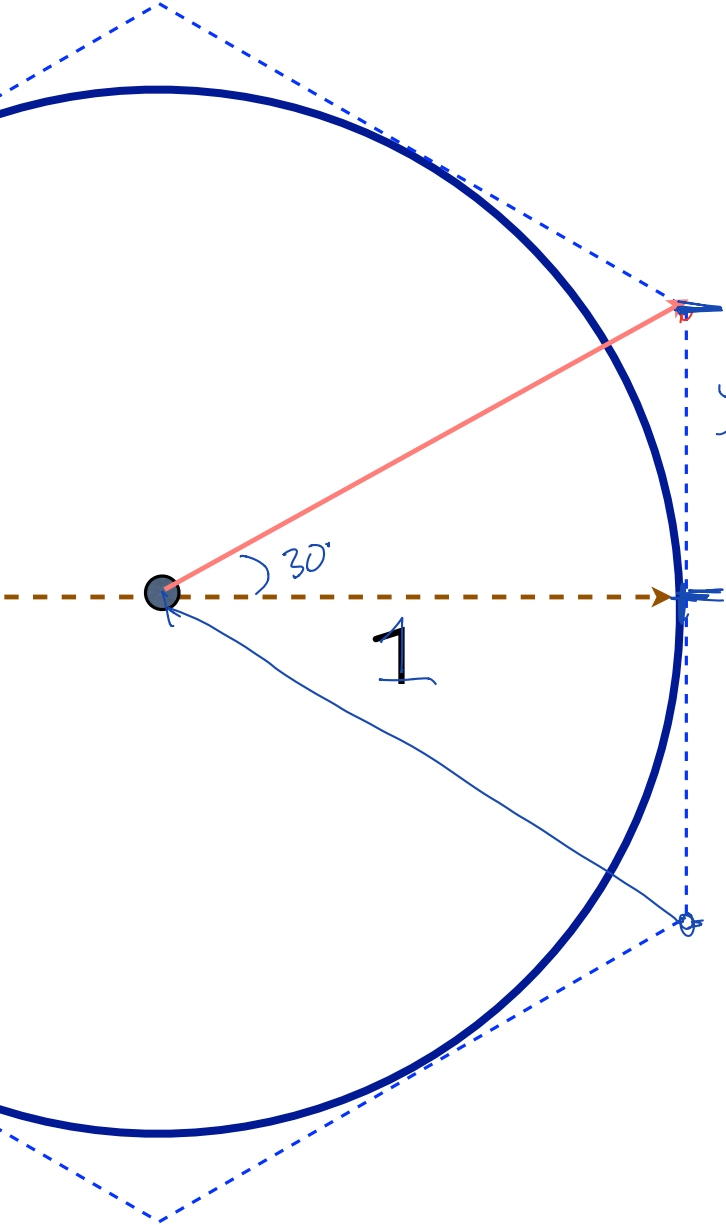




red perimeter $< \pi d$



6
 3
 red perimeter $< \pi d^2 <$ blue perimeter $4\sqrt{3}$ $2\sqrt{3}$



$$\tan(30) = \frac{y}{1} = \frac{1}{\sqrt{3}}$$

$$y = \frac{1}{\sqrt{3}}$$

entire blue perimeter =

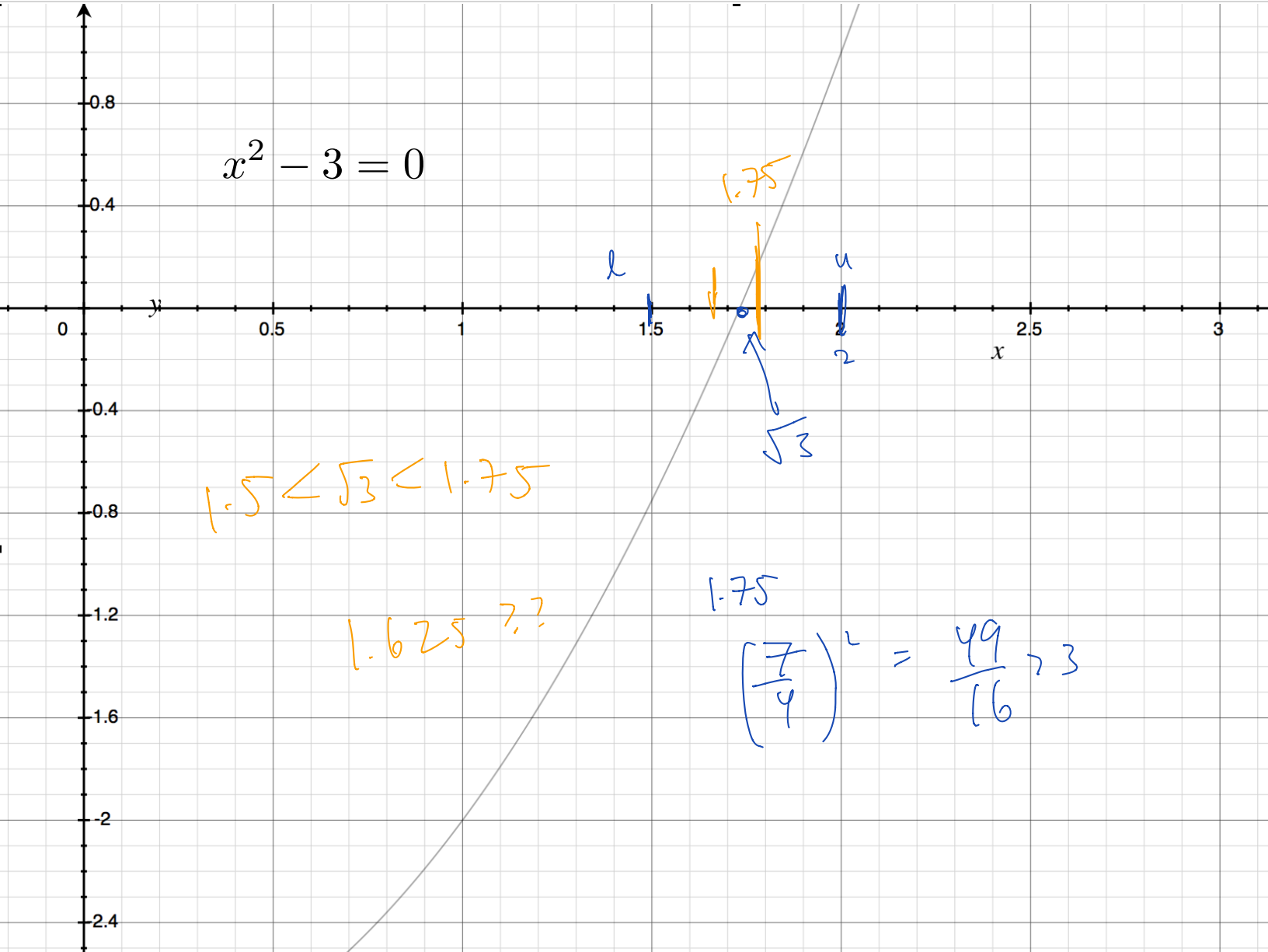
$$\frac{12}{\sqrt{3}} = \underline{\underline{4 \cdot \sqrt{3}}}$$

But what is

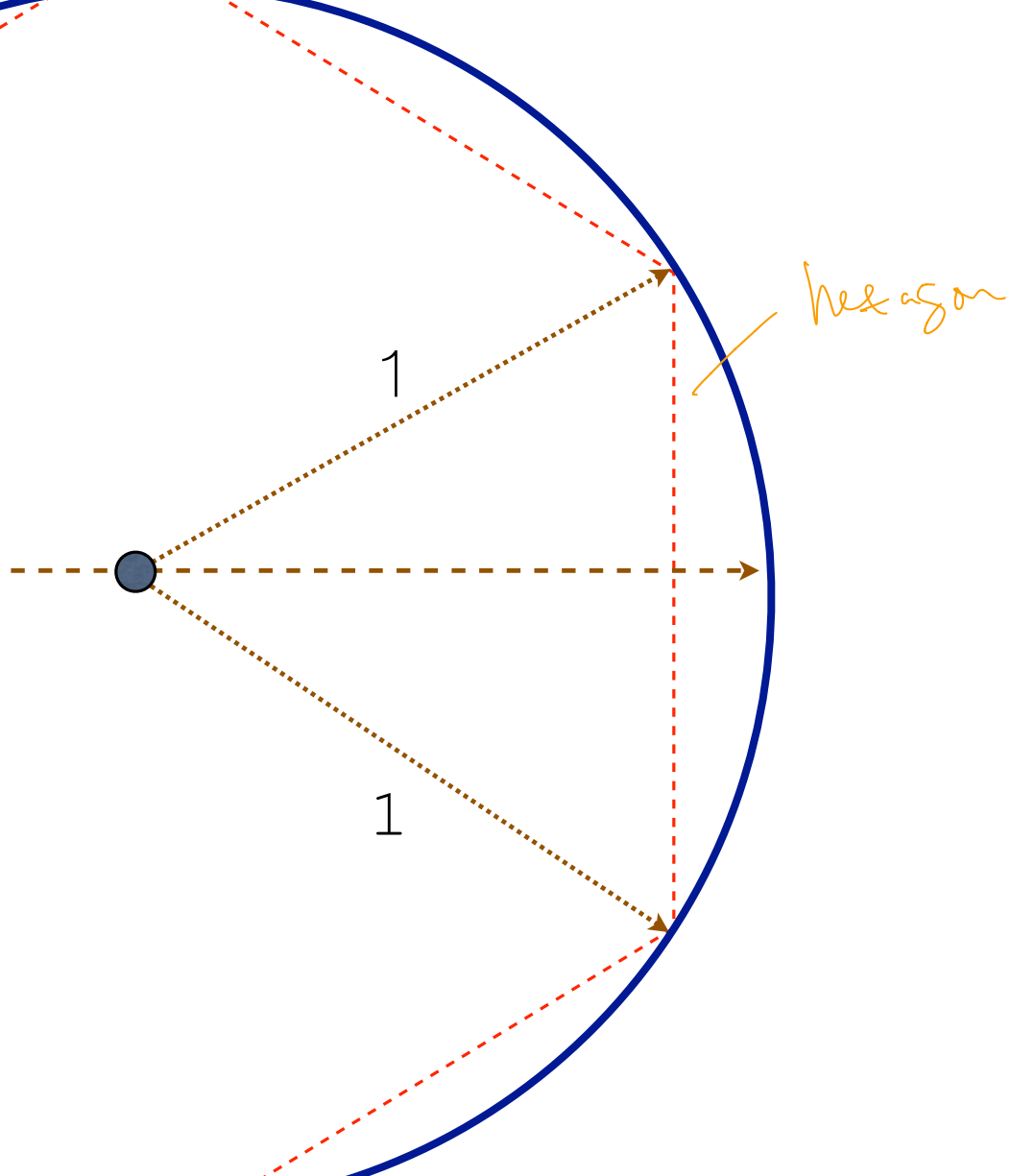
$$\sqrt{3}$$

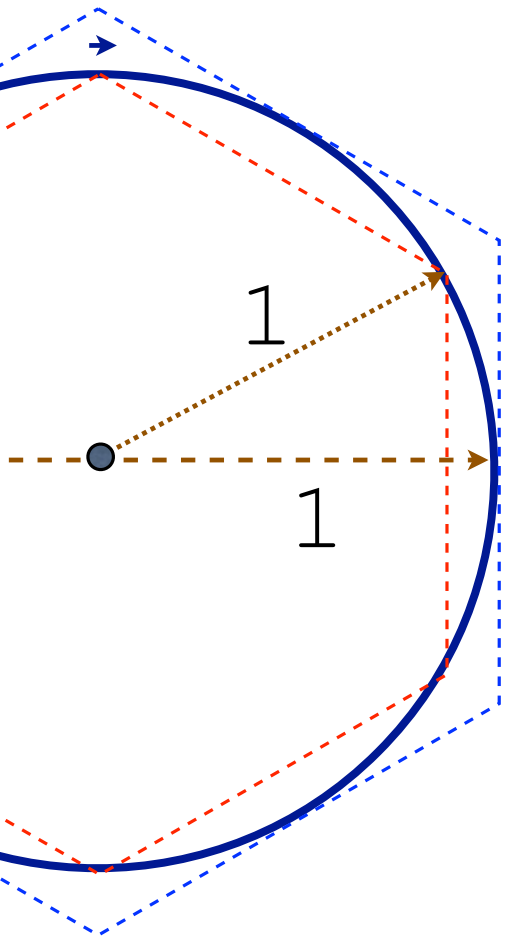
Theme 1 : reduce
the main problem
to a simpler one

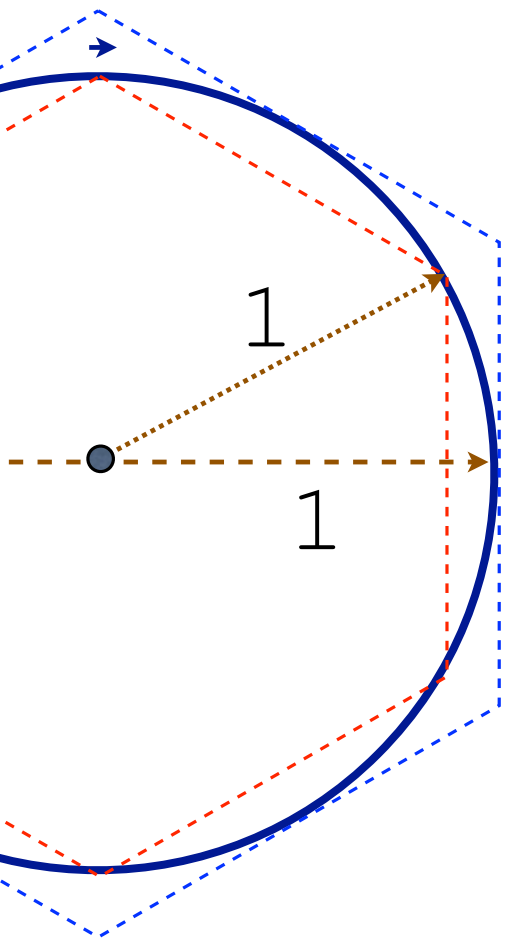
$$x^2 - 3 = 0$$



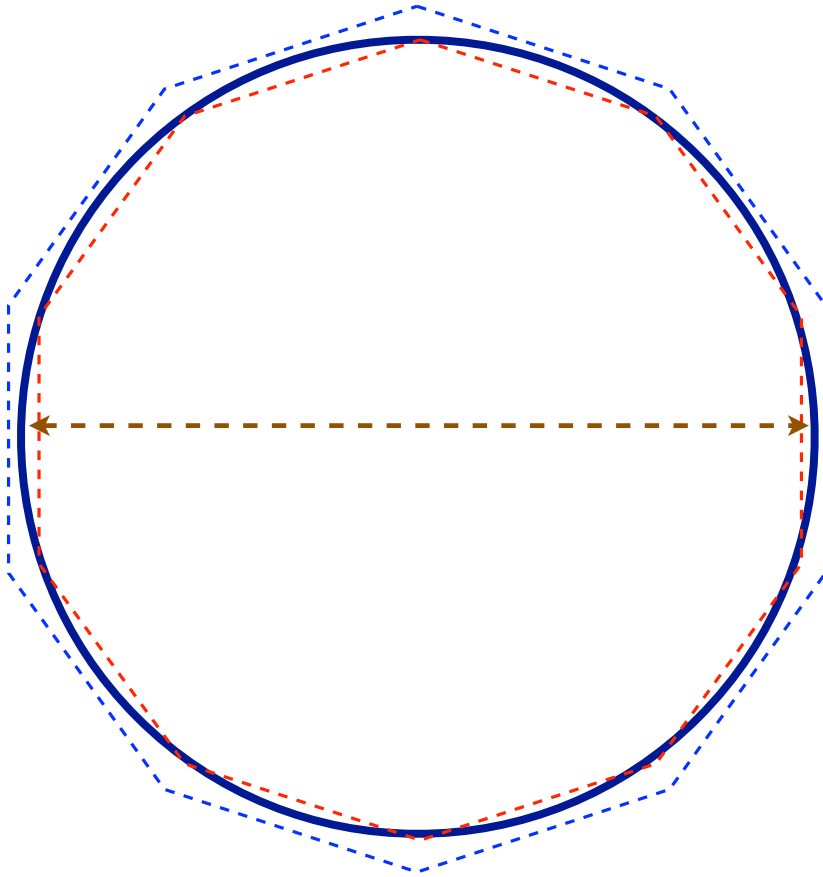
$$\frac{265}{153} \approx \sqrt{3}$$







red perimeter $< \pi d <$ blue perimeter



$$\begin{array}{r} 3.1428 \\ \hline \end{array} > \pi > \begin{array}{r} 3.1408 \\ \hline \end{array}$$

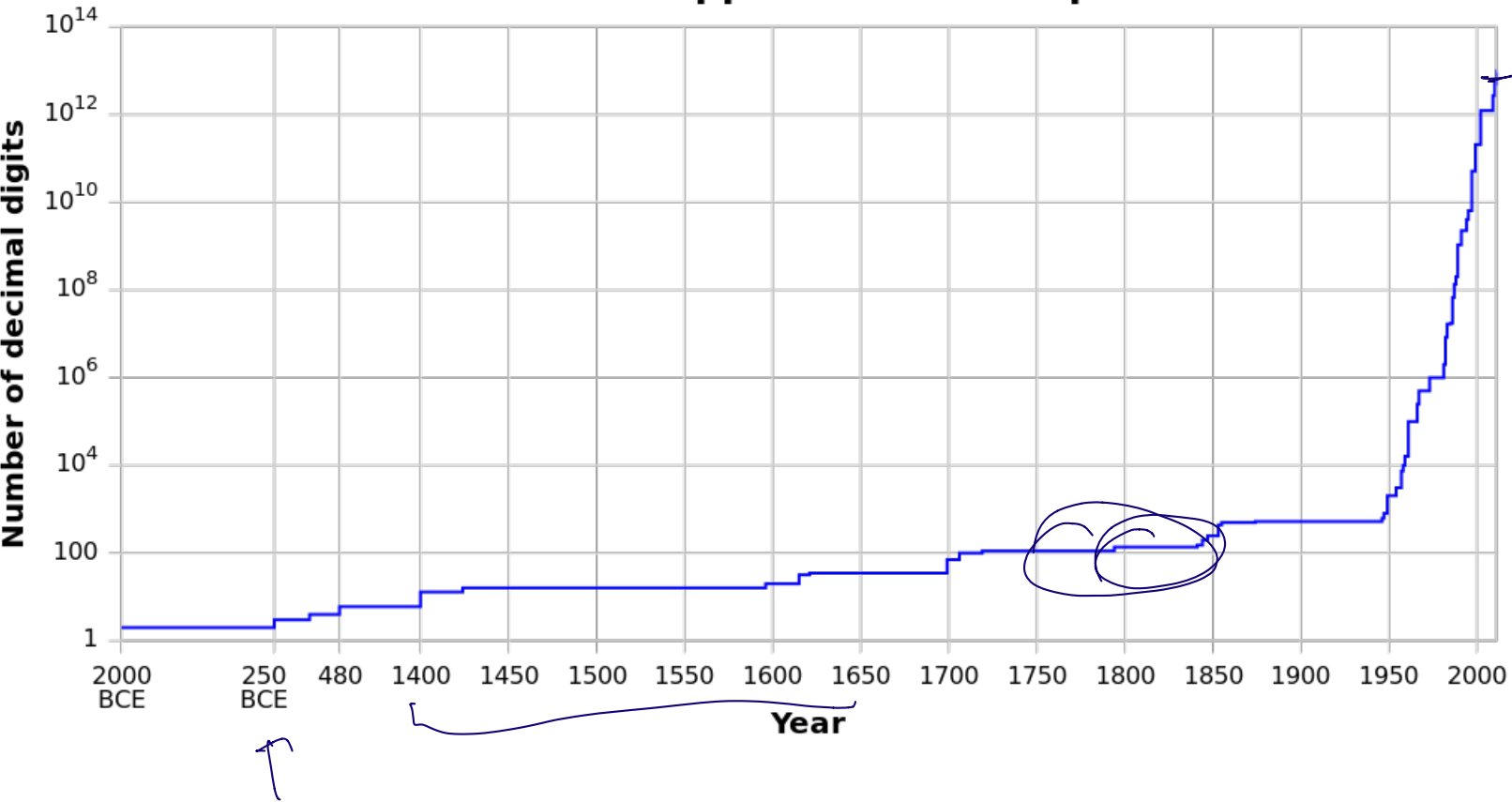
Using 96-gon,
Archimedes

how to analyze this approach?

efficiency analysis.

— accuracy versus work
trade off

Record approximations of pi



Theme3: new
insights lead to
improved
efficiency



$$\pi = \frac{9801}{\sqrt{8}} \left(\sum_{n=0}^{\infty} \frac{(4n)!(1103 + 26390n)}{(n!)^4 396^{4n}} \right)^{-1}$$

$$\pi = \frac{9801}{\sqrt{8}} \left(\sum_{n=0}^{\infty} \frac{(4n)!(1103 + 26390n)}{(n!)^4 396^{4n}} \right)^{-1}$$

$$n=0$$

$$\frac{(0!)(1103)}{0! \cdot 396^0}$$

$$\pi = \frac{9801}{\sqrt{8}} \cdot \left(\frac{1}{1103} \right)$$

$$\pi = \frac{9801}{\sqrt{8}} \left(\sum_{n=0}^{\infty} \frac{(4n)!(1103 + 26390n)}{(n!)^4 396^{4n}} \right)^{-1}$$

$$n=0$$

$$\pi \approx_0 \frac{9801}{\sqrt{8}} [1103]^{-1}$$

3.14159273001330576017

$$\pi = \frac{9801}{\sqrt{8}} \left(\sum_{n=0}^{\infty} \frac{(4n)!(1103 + 26390n)}{(n!)^4 396^{4n}} \right)^{-1}$$

n=1

$$\pi = \frac{9801}{\sqrt{8}} \left(\sum_{n=0}^{\infty} \frac{(4n)!(1103 + 26390n)}{(n!)^4 396^{4n}} \right)^{-1}$$

$$n=1$$

$$\pi \approx_1 \frac{9801}{\sqrt{8}} \left[1103 + \frac{24 \cdot 27493}{396^4} \right]^{-1}$$

3.14159265358979387799890582630

benefits?

much faster convergence for
the amount of work.

good algorithms touch
every aspect of our
lives

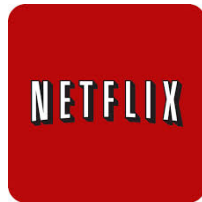
FedEx



Google

TESLA MOTORS

amazon



at&t



good algorithms
defend freedom



what skills do you
need for this course?

- precision

- creativity

in·ge·nu·i·ty

how to learn
in this class

no cookbook

develop
general problem
solving
skills

understand
known
techniques



work with your
peers

work with your
peers

but do not copy

A blue wavy underline consisting of two parallel lines that curve slightly upwards in the middle.

https://

shelat.khoury

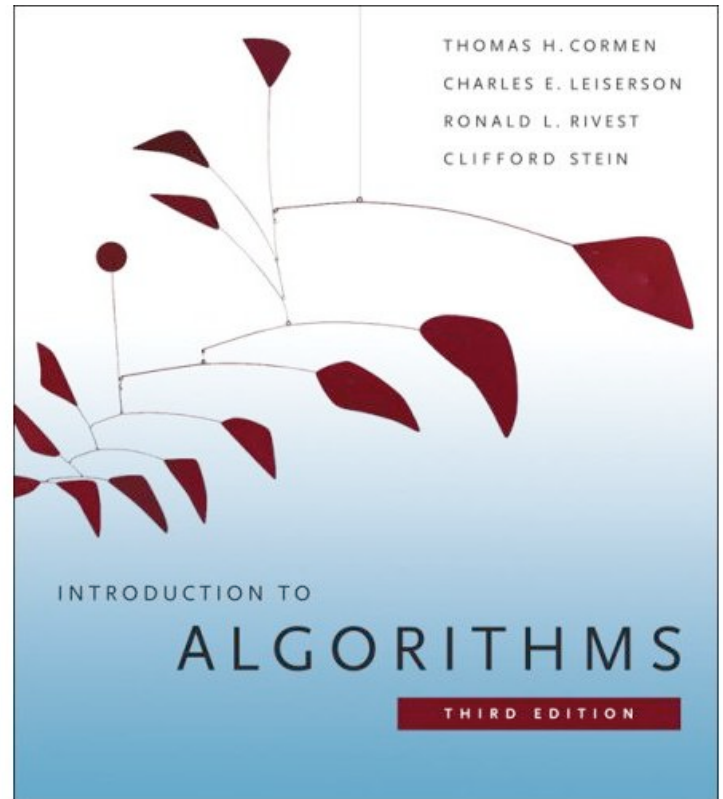
.northeastern

.edu/22s-5800

INTRODUCTION TO
ALGORITHMS

SECOND EDITION

THOMAS H. CORMEN
CHARLES E. LEISERSON
RONALD L. RIVEST
CLIFFORD STEIN



Availability

- This Month 53
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- 2-3 months 1

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All Topics

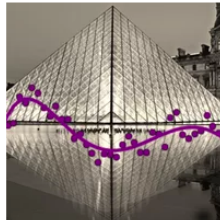
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You searched for algorithms. 118 matches



Algorithms, Part I

Princeton University



Approximation Algorithms Part I

École normale supérieure

LATEX



guide to latex



Search

About 41,900,000 results

Web

[\[PDF\] The Not So Short Introduction to LaTeX - Tobi Oetiker - Oetiker+ ...](#)

tobi.oetiker.ch/lshort/lshort.pdf

Images

a **LATEX** installation is available, ready to use. Information on how to access the local **LATEX** installation should be provided in the Local **Guide** [5]. If you.

Videos

News

[LaTeX - Wikibooks, open books for an open world](#)

en.wikibooks.org/wiki/LaTeX - Cached

Shopping

This is a **guide** to the **LaTeX** markup language. It is intended to form a useful resource for everybody from new users who wish to learn, to old hands who need a ...

More

[LaTeX/Mathematics - LaTeX/Document Structure - LaTeX/Text Formatting - Links](#)

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[\[PDF\] Short Math Guide for LaTeX - FTP Directory Listing - American ...](#)

ftp://ftp.ams.org/pub/tex/doc/amsmath/short-math-guide.pdf

Short Math **Guide** for **LATEX**. Michael Downes. American Mathematical Society. Version 1.09 (2002-03-22), currently available at.

[Guide to LaTeX \(4th Edition\): Helmut Kopka, Patrick W. Daly ...](#)

www.amazon.com/Guide-LaTeX-Edition-Helmut.../0321173856 - Cached

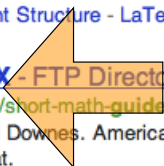
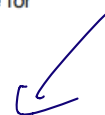
Guide to LaTeX (4th Edition) [Helmut Kopka, Patrick W. Daly] on Amazon.com. * FREE* super saver shipping on qualifying offers. Published Nov 25, 2003 by ...

[\[PDF\] A Beginner's Guide to LATEX - Princeton University](#)

www.cs.princeton.edu/courses/archive/spr10/cos433/Latex/latex-guide.pdf - Cached - Similar

A Beginner's **Guide to LATEX**. David Xiao dxiao@cs.princeton.edu. September 12, 2005. 1 Introduction. LATEX is the standard mathematical typesetting ...

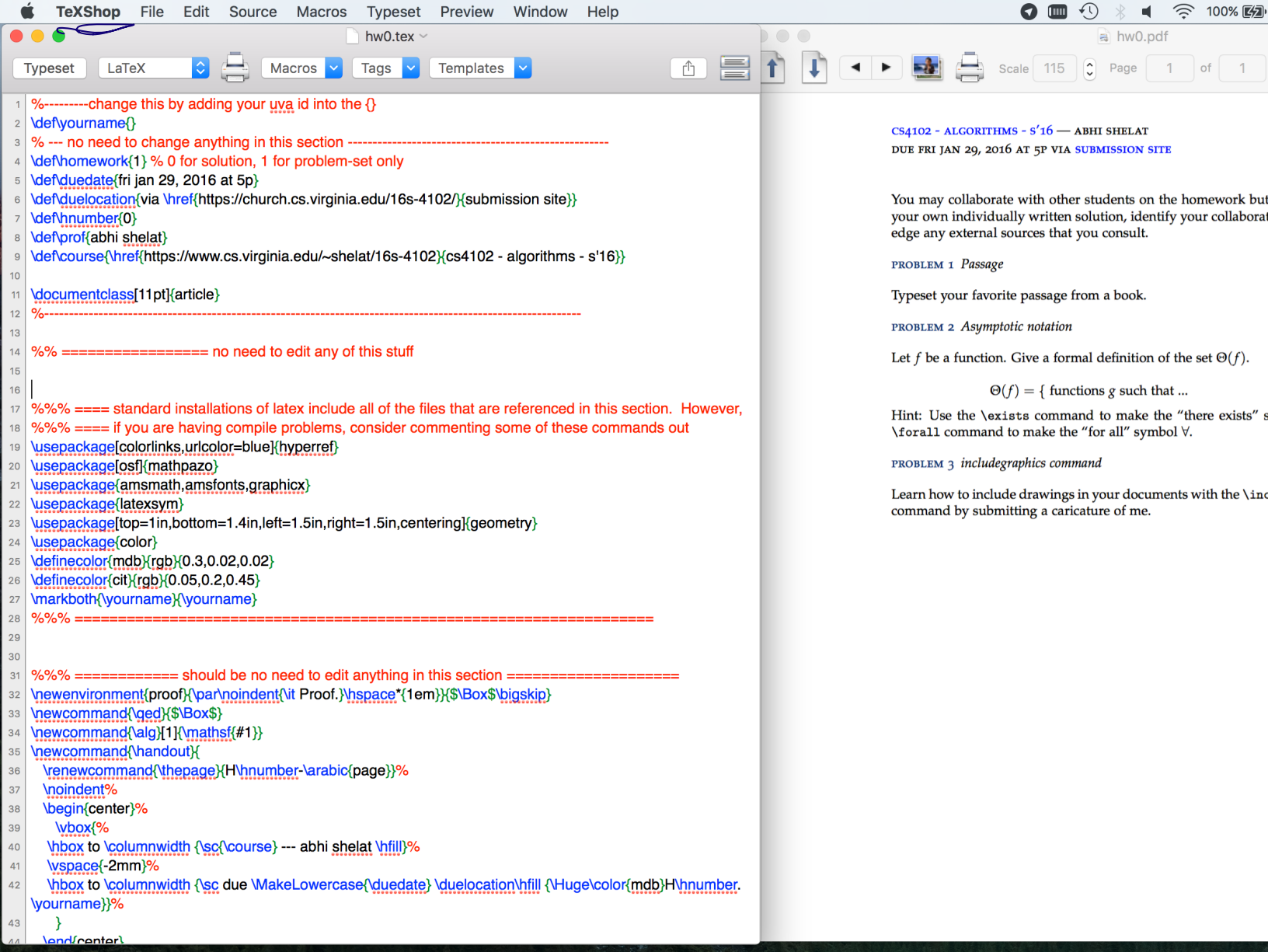
[LaTeX documentation](#)



The Not So Short Introduction to L^AT_EX 2_ε

Or E_T_X 2_ε in *157* minutes

by Tobias Oetiker



```

1 %-----change this by adding your uva id into the {}
2 \def\yourname{}
3 % --- no need to change anything in this section -----
4 \def\homework{1} % 0 for solution, 1 for problem-set only
5 \def\duedate{fri jan 29, 2016 at 5p}
6 \def\duelocation{via \href{https://church.cs.virginia.edu/16s-4102/}{submission site}}
7 \def\hnumber{0}
8 \def\prof{abhi shelat}
9 \def\course{\href{https://www.cs.virginia.edu/~shelat/16s-4102}{cs4102 - algorithms - s'16}}
10
11 \documentclass[11pt]{article}
12 %-----
13
14 %% ===== no need to edit any of this stuff
15
16 |
17 %%% ==== standard installations of latex include all of the files that are referenced in this section. However,
18 %%% ==== if you are having compile problems, consider commenting some of these commands out
19 \usepackage{colorlinks,urlcolor=blue}{hyperref}
20 \usepackage{osf}{mathpazo}
21 \usepackage{amsmath,amsfonts,graphicx}
22 \usepackage{latexsym}
23 \usepackage{top=1in,bottom=1.4in,left=1.5in,right=1.5in,centering}{geometry}
24 \usepackage{color}
25 \definecolor{mdb}{rgb}{0.3,0.02,0.02}
26 \definecolor{cit}{rgb}{0.05,0.2,0.45}
27 \markboth{\yourname}{\yourname}
28 %%% =====
29
30
31 %%% ===== should be no need to edit anything in this section =====
32 \newenvironment{proof}{\par\noindent\it Proof.}\hspace*{1em}}{\Box$\bigskip}
33 \newcommand{\qed}{\Box$}
34 \newcommand{\alg}[1]{\maths{#1}}
35 \newcommand{\hayout}{
36   \renewcommand{\thepage}{H\hnumber-\arabic{page}}}%
37   \noindent%
38   \begin{center}%
39     \vbox{%
40       \hbox to \columnwidth {\sc\course} --- abhi shelat \hfill}%
41       \vspace{-2mm}%
42       \hbox to \columnwidth {\sc due \MakeLowercase{\duedate} \duelocation\hfill {\Huge\color{mdb}H\hnumber.
43 \yourname}}}%
44   \end{center}

```

CS4102 - ALGORITHMS - S'16 — ABHI SHELAT
DUE FRI JAN 29, 2016 AT 5P VIA [SUBMISSION SITE](#)

You may collaborate with other students on the homework but your own individually written solution, identify your collaborate edge any external sources that you consult.

PROBLEM 1 *Passage*

Typeset your favorite passage from a book.

PROBLEM 2 *Asymptotic notation*

Let f be a function. Give a formal definition of the set $\Theta(f)$.

$$\Theta(f) = \{ \text{functions } g \text{ such that } \dots$$

Hint: Use the `\exists` command to make the “there exists” symbol \exists .
Use the `\forall` command to make the “for all” symbol \forall .

PROBLEM 3 *includegraphics command*

Learn how to include drawings in your documents with the `\includegraphics` command by submitting a caricature of me.



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TeXShop

pages.uoregon.edu/koch/texshop/ ▾ University of Oregon ▾

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Obtaining TeXShop

Obtaining TeXShop. If you just want to upgrade to the latest ...

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Installing

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- Introduction
- Page headers and footers
 - What is \textsf{fancyhdr}
 - Simple use of \textsf{fancyhdr}
 - A simple example
- An example of two-sided printing
- Redefining \texttt{plain} style
- The default layout
- The scoop on \textsf{latex}'s marks
- Dictionary style headers
- Fancy layouts
 - Two book examples
 - Special page layout for float pages
 - Those blank pages
 - \textsf{N} of \textsf{M} style page numbers
 - Chapter or section related page numbers
 - When to change the headers and footers?
 - Headers and footers induced by the text
- Package for extra marks in \textsf{LaTeX}
- A movie
- Thumb indexes
- Float placement
- Multipage Floats
- Contact information

```

...
\extramarks{}(Continued on next page\ldots)
Some text that may or
\extramarks(Continued)
\end{verbatim}

\CmdIndex(extramarks)
Note that the \Cmd(extramarks) command must be close to the text, i.e no
empty lines (paragraph boundaries) should intervene. Otherwise the page may
be broken at that boundary and the extramarks would come on the wrong page.

There are two new marks that can be used in the page layout with this
package: If commands of the form
\verb|\extramarks(|$m_1$\verb|)(|$m_2$\verb|)| are given
\CmdIndex(firstxmark)
\CmdIndex(lastxmark)
\Cmd(firstxmark) gives you the first $m_1$ value and
\Cmd(lastxmark) gives you the last $m_2$ value
of the current page.
\CmdIndex(firstleftmark)
\CmdIndex(lastrightmark)
es you the \Cmd(firstleftmark) and \Cmd(lastrightmark)
at complement the standard \textsf{latex}/ marks.

he point that marks are the correct way to do this, let me
'solution' that will not work\footnote(Actually there is
ay but it requires two \textsf{latex}/ passes: you can put \Cmd(label)
before and after the text and compare the \Cmd(pageref)s.):

\begin{verbatim}
\lhead(Continued)

```

Greek small Letters

α β χ δ ε φ ψ γ η ι κ λ μ ν ο π ρ θ ρ σ ς τ υ ω ξ ψ ζ

- Go to
- Insert Label
- Insert Reference**
- Insert Reference to Page
- Properties...

```

-i-ni-tion would
Underfull \hbox (badness 5077) in paragraph at lines 1088--1095
/cmr10/be /cmctt10/\thechapter-\arabic{page} /cmr10/ but you can give this def-i-
ni-tion your-self af-ter the
[16] [17] [18] [19] [20] [21] [22] [23] [24]
No file fancyhdr.ind.
[25] (fancyhdr.aux)
LaTeX Warning: Label(s) may have changed. Rerun to get cross-references right.

```

Overleaf

The screenshot displays the Overleaf web interface for a project titled "5800 H0". The browser address bar shows the URL: <https://www.overleaf.com/project/61e6bd989113c8ca0a6ca65c>. The interface is divided into several sections:

- Top Bar:** Contains navigation icons and buttons for "Review", "Share", "Submit", "History", and "Chat".
- Editor Pane (Left):** Shows the LaTeX source code for "main.tex". The code includes:
 - Comments and macros for user identification: `\def\yourname{}`, `\def\homework{1}`, `\def\duedate{wed jan 26, 2022 at 11.59p}`, `\def\duelocation{via}`, and `\href{https://gradescope.com/courses/331917}{gradescope}`.
 - Document class and packages: `\documentclass[11pt]{article}`, `\usepackage{colorlinks,urlcolor=blue}{hyperref}`, `\usepackage[osf]{mathpazo}`, `\usepackage{amsmath,amsfonts,graphicx}`, `\usepackage{latexsym}`, `\usepackage[top=1in,bottom=1.4in,left=1.5in,right=1.5in,centering]{geometry}`, `\usepackage{color}`, `\definecolor{mdb}{rgb}{0.3,0.02,0.02}`, `\definecolor{cit}{rgb}{0.05,0.2,0.45}`, and `\markboth{\yourname}{\yourname}`.
 - Comments indicating sections that do not need editing.
- Rendered PDF (Right):** Shows the output of the compilation. The header includes "CS5800 ALGORITHMS S'22 — ABHI SHELAT" and "DUE WED JAN 26, 2022 AT 11:59P VIA GRADESCOPE". The main content features the name "Ho." in a large serif font, followed by instructions on collaboration and submission, and a section titled "PROBLEM 1 Passage" with the instruction "Typeset your favorite passage from a book."

Submitting HW

[gradescope](#)

Honor Policy



I, _____, do hereby certify on my honor that during this course,

1. I shall write my answers entirely by myself, and neither share nor request text, code, or drawings.
2. I will not give or derive assistance from any unauthorized sources or the web.

today

midterm
Feb 23

final
~~Dec~~
May



today

midterm
Feb 23

final
May 4

dynamic

graphs



div & conq

greedy

Lp,np,randomized

counting

First example of an
algorithmic pattern based
on I1 and I2

1

stand

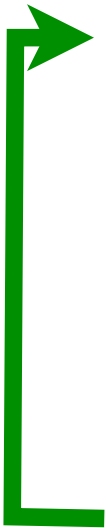
1 stand

2 set your “number” to one

- 1 stand
- 2 set your “number” to one
- 3 greet your neighbor (pause if no partner)

- 1 stand
- 2 set your “number” to one
- 3 greet your neighbor (pause if no partner)
- 4 if you are older, give “number” and sit
if you are younger, add “numbers”

- 1 stand
- 2 set your “number” to one
- 3 greet your neighbor (pause if no partner)
- 4 if you are older, give “number” and sit
if you are younger, add “numbers”
- 5 if you are standing & you have a neighbor,
goto 3



1

stand

2

set

3

greet

4

sit/add

5

repeat

lets analyze this alg

Our model of computation

Basic op: 1 unit

Set, Greet, add, compare, sit

Simplify: in each round, every standing person can do 1 op

Lets count # of rounds until we finish

1

stand

2

set

3

greet

4

sit/add

5

repeat



how fast does it work:

1

stand

2

set

3

greet

4

sit/add

5

repeat



how fast does it work:

$T(n)$

rounds to finish in a room with n people

1

stand

1

2

set

1

3

greet

1

4

sit/add

5

repeat



Simple case: 1 person

$$T(1) = 3$$

1

stand

|

2

set

|

3

greet

||

4

sit/add

|

5

repeat



Simple case: 2 people

$$T(2) = 5$$

1

stand



2

set



3

greet



4

sit/add



5

repeat



$$T(4) = 4 + \underline{\underline{T(2)}}$$

1

stand

2

set

3

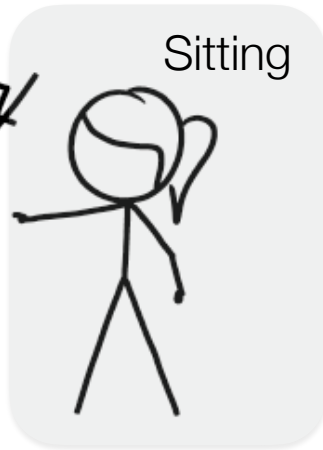
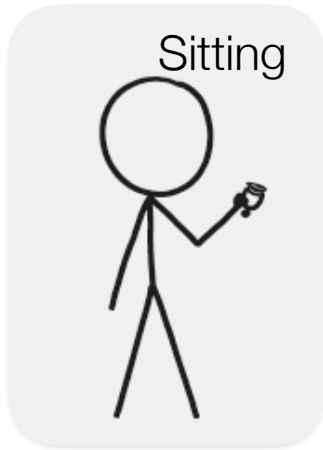
greet

4

sit/add

5

repeat



After step 4

$$T(4) =$$

1 2
stand set

3 4 5
greet sit/add repeat

These steps only
happen once.

What about these?

I1: Approx is OK

1 2
stand set

3 4 5
greet sit/add repeat

how fast does it work:

$$T(\underline{n}) = 1 + 1 + \underbrace{T(\lceil n/2 \rceil)}$$



how fast does it work:

$$T(n) = 1 + 1 + T(\lceil n/2 \rceil)$$

$$T(1) = 3$$

This is a recurrence

$$T(n) = T(\lceil n/2 \rceil) + 2$$

$$T(1) = 3$$

solve a simpler case when n is a power of 2.

$$T(2^k) = 2 + T(2^{k-1})$$

$$T(2^k) = 2 + T(2^{k-1})$$

$$\begin{aligned} T(2^k) &= 2 + T(2^{k-1}) \\ &= 2 + 2 + T(2^{k-2}) \end{aligned}$$

$$\begin{aligned} T(2^k) &= 2 + T(2^{k-1}) \\ &= 2 + 2 + T(2^{k-2}) \\ &= \overbrace{2 + 2 + \cdots + 2}^k + T(2^0) \end{aligned}$$

$$\begin{aligned} T(2^k) &= 2 + T(2^{k-1}) \\ &= 2 + 2 + T(2^{k-2}) \\ &= \overbrace{2 + 2 + \cdots + 2}^k + T(2^0) \\ &= 2k + T(1) \end{aligned}$$

“intuition here”

$$T(2^k) = 2 + T(2^{k-1})$$

Other cases?

“intuition here”

$$\begin{aligned}T(2^k) &= 2 + T(2^{k-1}) \\ &= 2 + 2 + T(2^{k-2})\end{aligned}$$

Other cases?

“intuition here”

$$\begin{aligned}T(2^k) &= 2 + T(2^{k-1}) \\ &= 2 + 2 + T(2^{k-2}) \\ &= \overbrace{2 + 2 + \cdots + 2}^k + T(2^0)\end{aligned}$$

Other cases?

“intuition here”

$$\begin{aligned}T(2^k) &= 2 + T(2^{k-1}) \\ &= 2 + 2 + T(2^{k-2}) \\ &= \overbrace{2 + 2 + \cdots + 2}^k + T(2^0) \\ &= 2k + T(1)\end{aligned}$$

Other cases?

Idea1 : It is OK
to approximate

Asymptotic notation

$O(g)$

This notation represents a set

Asymptotic notation

$O(g)$

Set of functions that are at most within const of g for large n

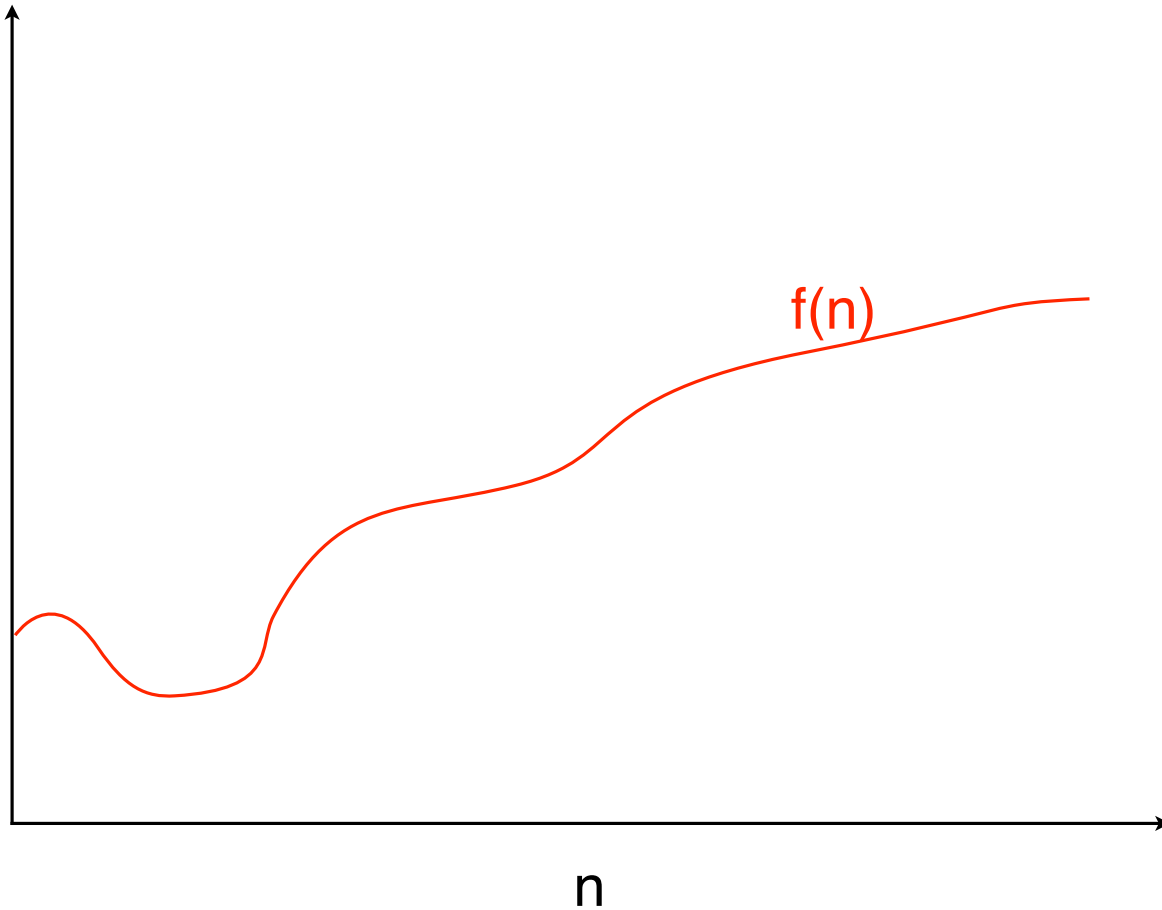
Asymptotic notation

$O(g)$ at most within const of g for large n

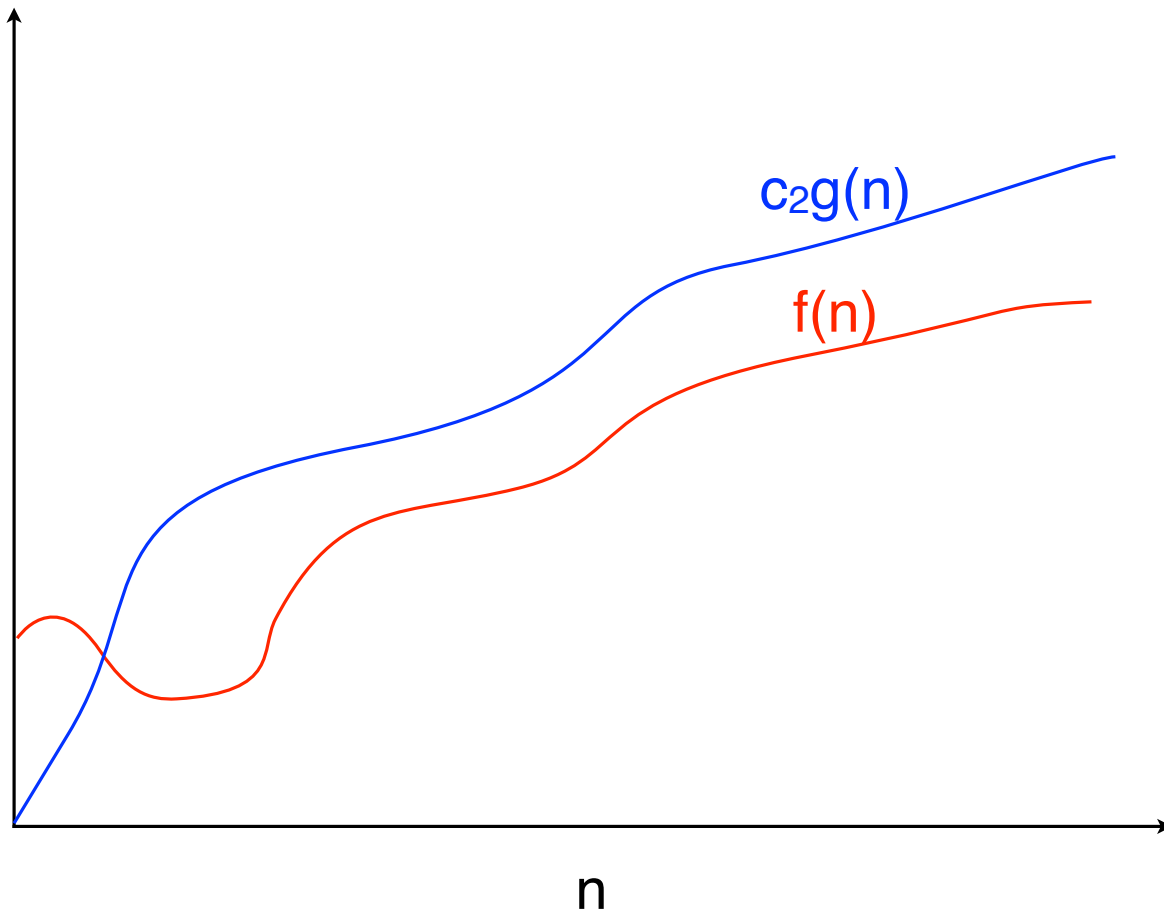
$\Omega(g)$ at least within const of g for large n

$\Theta(g)$ within a const of g for large n

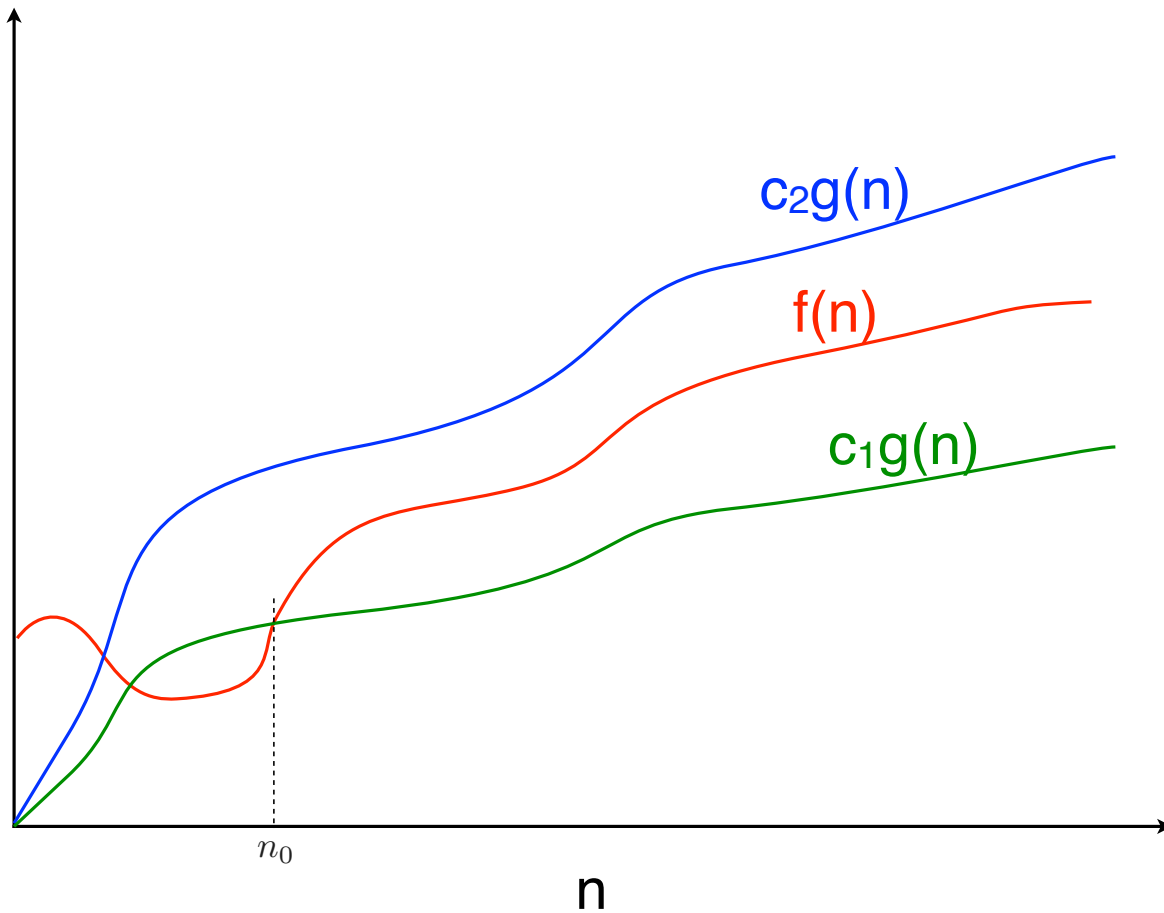
Omega sandwich



Omega sandwich



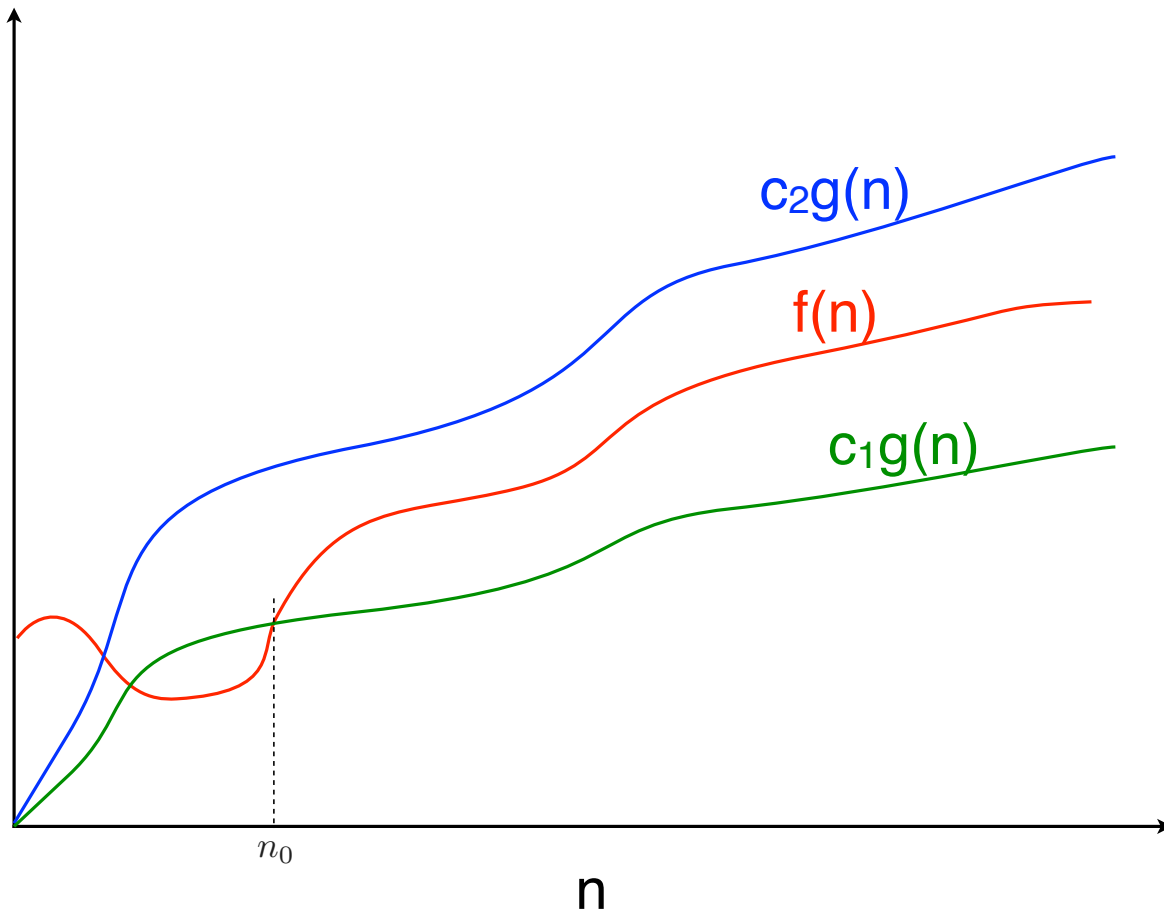
Omega sandwich



$$f(n) = O(g(n))$$

$$f(n) = \Omega(g(n))$$

Omega sandwich



$$f(n) = O(g(n))$$

$$f(n) = \Theta(g(n))$$

$$f(n) = \Omega(g(n))$$

“intuition here”

$$T(2^k) = 2 + T(2^{k-1})$$

$$= 2 + 2 + T(2^{k-2})$$

$$= \overbrace{2 + 2 + \cdots + 2}^k + T(2^0)$$

$$= 2k + T(1) = O(\log(2^k))$$

“intuition here”

$$T(2^k) = 2 + T(2^{k-1})$$

$$= 2 + 2 + T(2^{k-2})$$

$$= \overbrace{2 + 2 + \cdots + 2}^k + T(2^0)$$

$$= 2k + T(1) = O(\log(2^k))$$

$$\forall 0 < n < m, T(n) \leq T(m)$$

$$T(m) \leq T(2^{\lceil \log(m) \rceil}) = 2 \lceil \log(m) \rceil + 2$$

“intuition here”

$$T(2^k) = 2 + T(2^{k-1})$$

$$= 2 + 2 + T(2^{k-2})$$

$$= \overbrace{2 + 2 + \cdots + 2}^k + T(2^0)$$

$$= 2k + T(1) = O(\log(2^k))$$

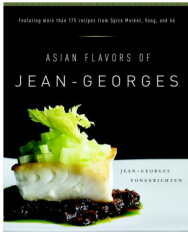
$$\forall 0 < n < m, T(n) \leq T(m)$$

$$T(m) \leq T(2^{\lceil \log(m) \rceil}) = 2 \lceil \log(m) \rceil + 2$$

$$T(m) = \Omega(\log(m))$$

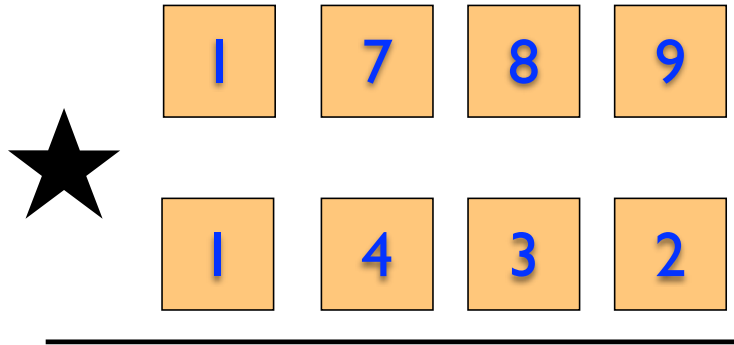
$$= \Theta(\log(m))$$

How to solve recurrence relations

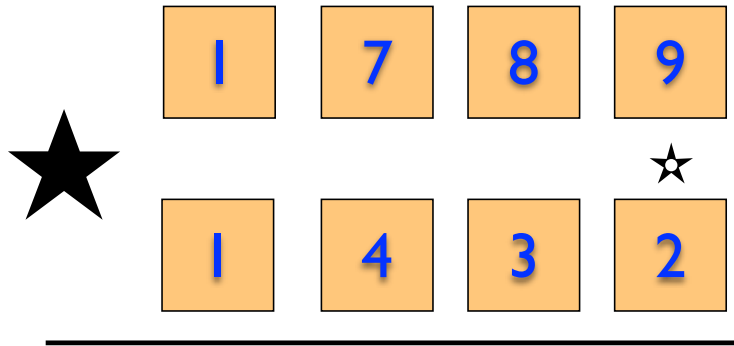


Multiplication

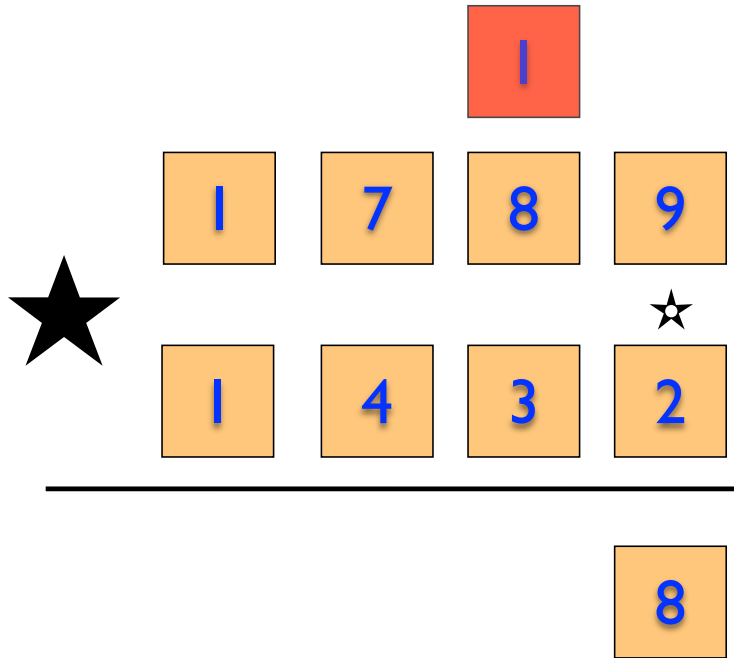
How much work does it take?



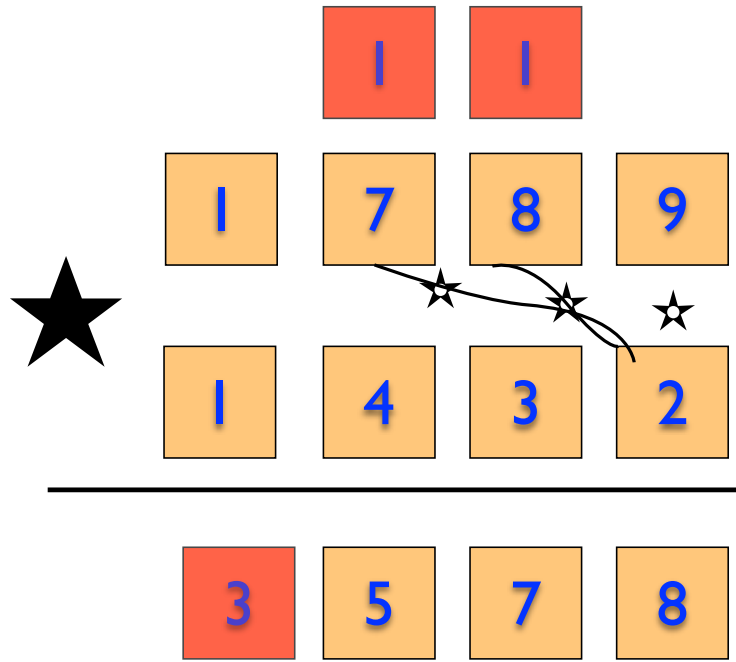
$$(n-1)(n+1) +$$



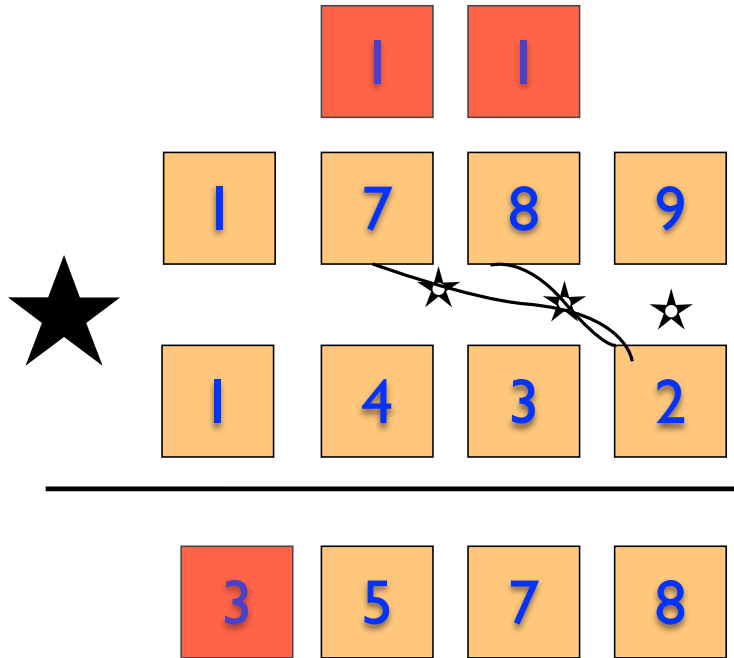
$$(n-1)(n+1) +$$



$$(n-1)(n+1) +$$

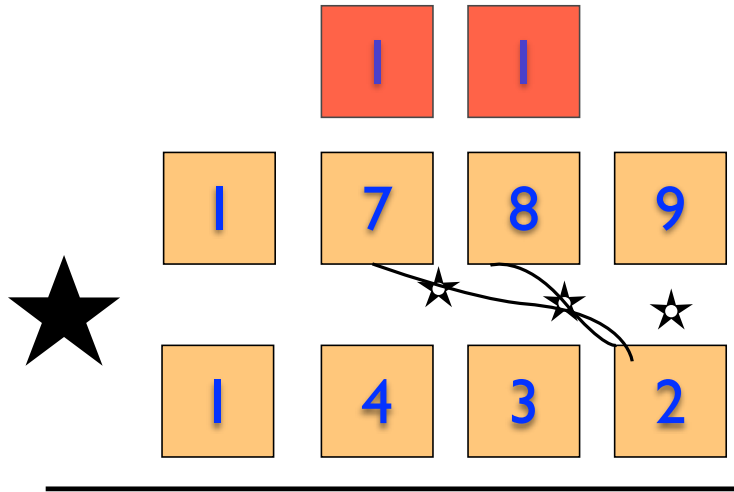


$$(n-1)(n+1) +$$

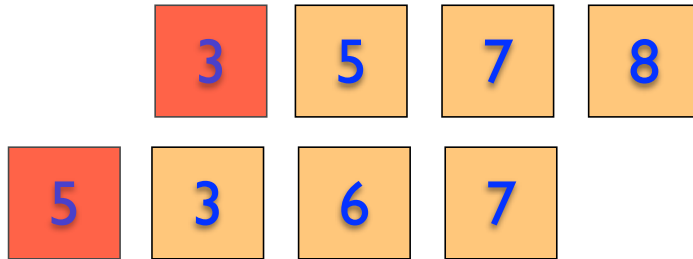


$$(n-1)(n+1) +$$

$$n \star \quad n-1 +$$

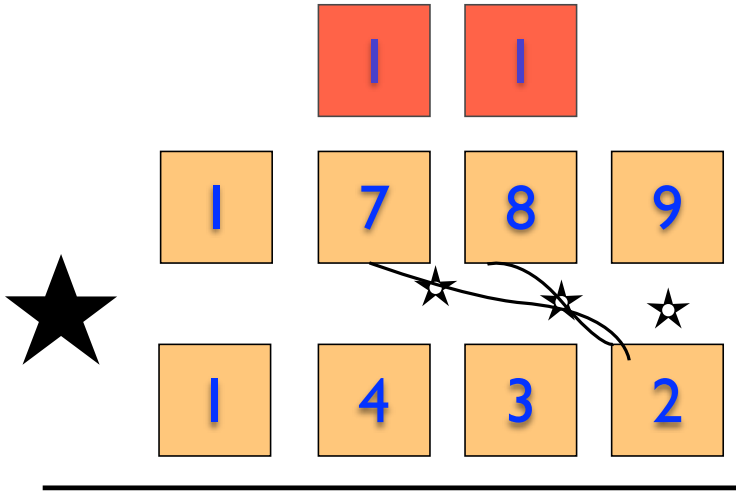


$$(n-1)(n+1) +$$

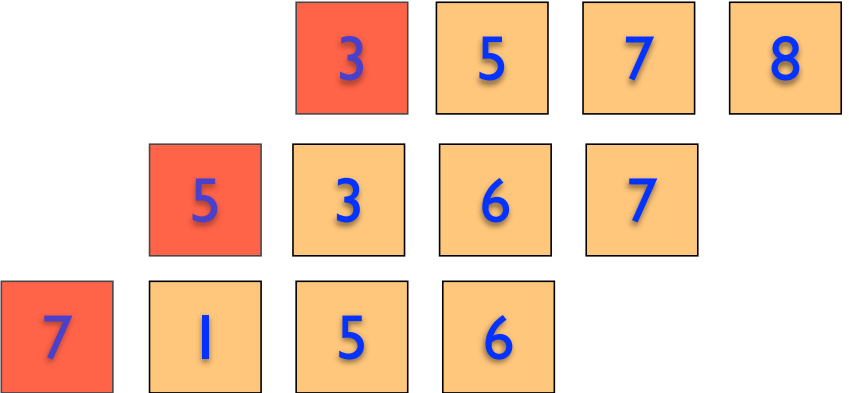


$$n \star \quad n-1 +$$

$$n \star \quad n-1 +$$



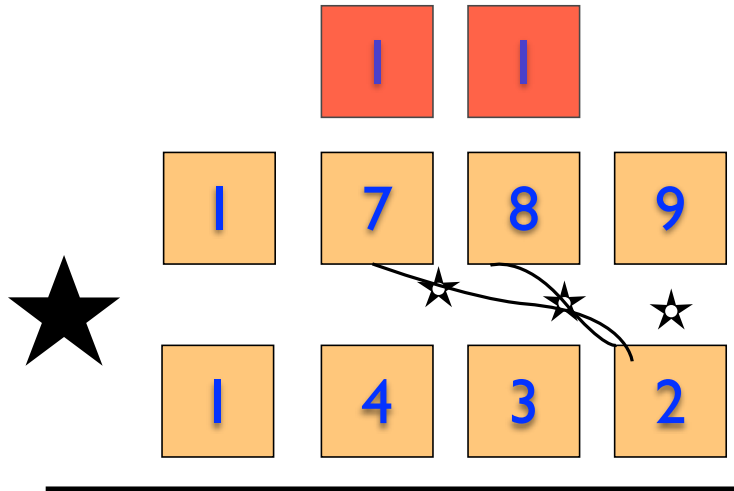
$$(n-1)(n+1) +$$



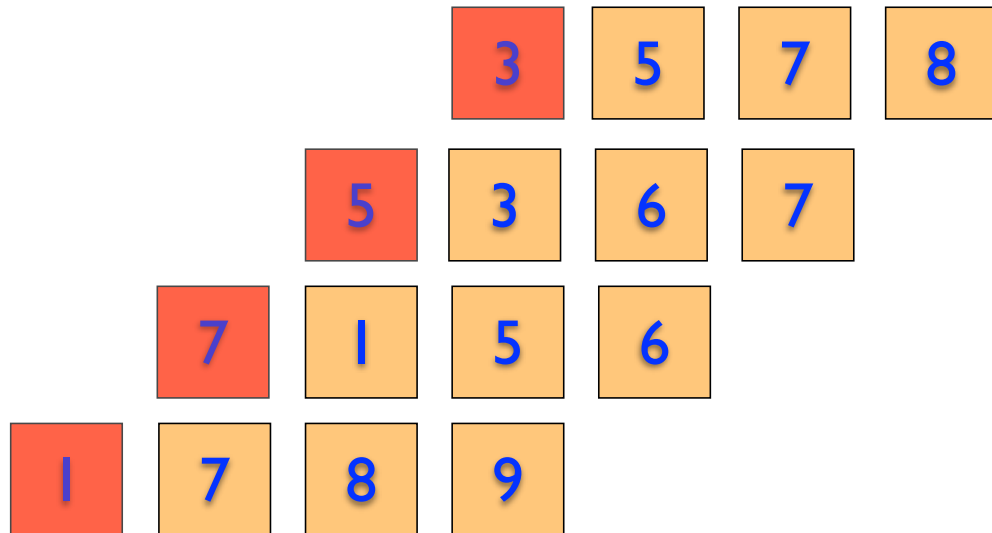
$$n \star \quad n-1 +$$

$$n \star \quad n-1 +$$

$$n \star \quad n-1 +$$



$$(n-1)(n+1) +$$

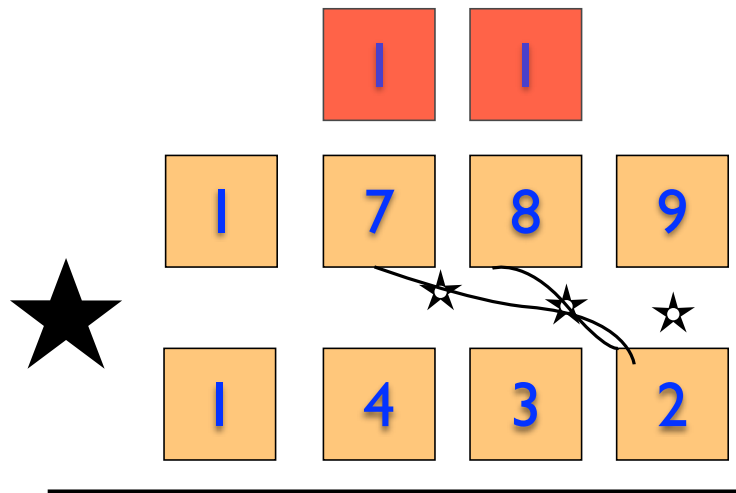


$$n \star \quad n-1 +$$

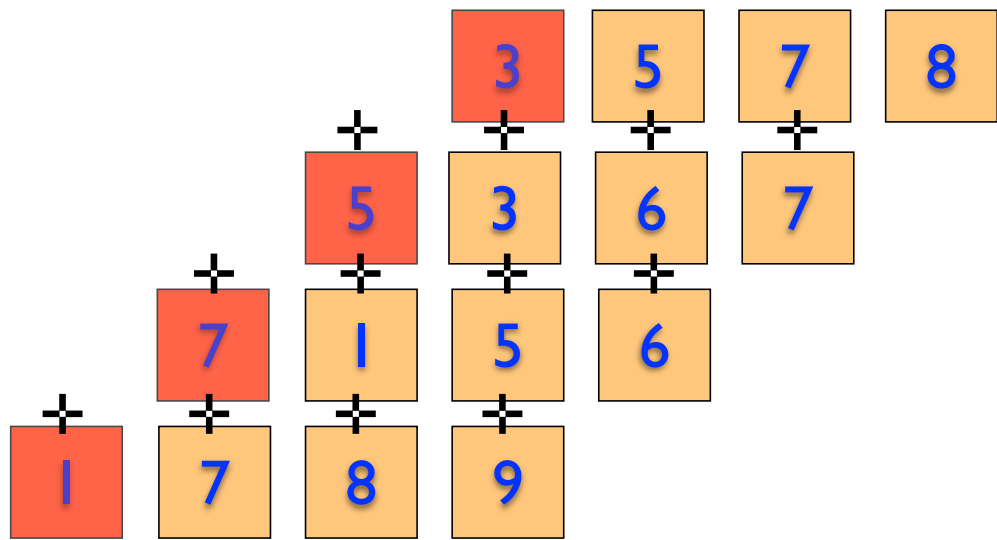
$$n \star \quad n-1 +$$

$$n \star \quad n-1 +$$

$$n \star \quad n-1 +$$



$$(n-1)(n+1) +$$



$$n \star \quad n-1 +$$

$$n \star \quad n-1 +$$

$$n \star \quad n-1 +$$

$$n \star \quad n-1 +$$

Theme 1

A first attempt...



1 7 8 9



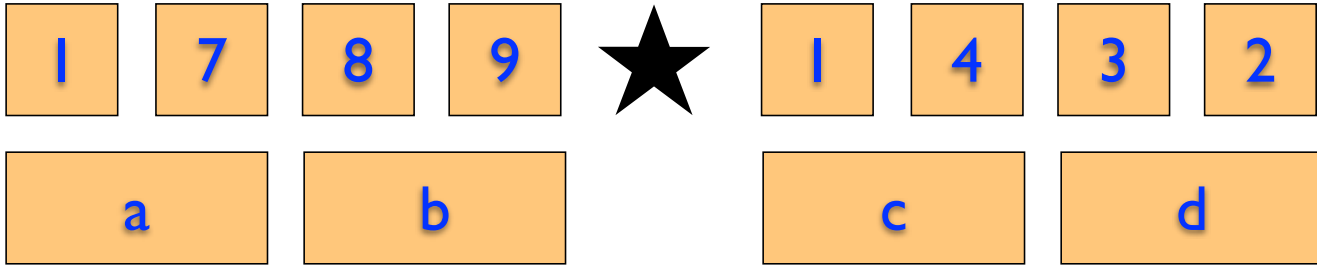
1 4 3 2

a

b

c

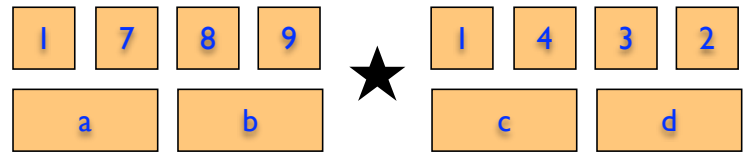
d



$$ac100^2 + (ad + bc)100 + bd$$

n-digit inputs

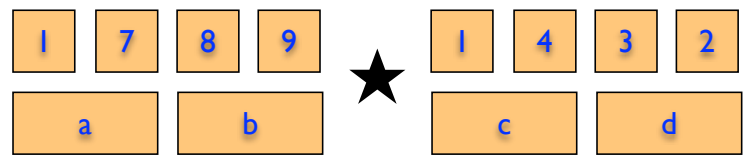
Mult(ab, cd)



$$ac100^2 + (ad + bc)100 + bd$$

Base case: return $b*d$ if inputs are 1-digit

Mult(ab, cd)



$$ac100^2 + (ad + bc)100 + bd$$

Base case: return $b*d$ if inputs are 1-digit

Compute $x = \text{Mult}(a,c)$

Compute $y = \text{Mult}(a,d)$

Compute $z = \text{Mult}(b,c)$

Compute $w = \text{Mult}(b,d)$

Return $r = x*100^2 + (y+z)100 + w$

$$T(n) = 4T(n/2) + 3O(n)$$



calculations: